

Optimal Power Control Over Multiple Time-Scale Fading Channels With Service Outage Constraints

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Abstract—This paper considers the power-control problem for a fading channel in an information-theoretic framework. We derive power-control schemes to optimize ergodic capacity, outage capacity, and capacity with a service outage constraint. The novelty in the paper lies in the use of a two-time-scale fading process and its implications for the channel-state information available at the transmitter.

Index Terms—Ergodic capacity, fading channels, outage capacity, power control, service outage.

I. INTRODUCTION

DETERMINING the information-theoretic capacity of fading wireless channels has been an important area of research over the past decade. Interest in this area has been spurred on by the tremendous growth in mobile and wireless communications networks, from cellular systems to wireless local area networks. Importantly, many recent information-theoretic results are impacting the design of next-generation wireless networks, which will include techniques such as adaptive modulation and coding and channel-based scheduling to exploit multiuser diversity. While numerous problems arise in multiuser environments such as multiple-access channels and broadcast channels, there are still many interesting questions to be answered about the capacity of single-user fading channels. This paper addresses one such question.

There are various notions of *capacity* for single-user fading channels, the main ones being *ergodic capacity* [1], *delay-limited capacity* [2], and *capacity versus outage probability* [3], [4]. Ergodic capacity is a capacity that can be achieved by averaging over all states of an ergodic fading channel, and thus, is suitable for non-real-time traffic applications. Delay-limited capacity and capacity versus outage concepts are useful for constant-rate real-time traffic applications. An excellent survey of various information-theoretic notions for fading channels is given in [5].

Another important issue in studying the capacity of fading channels is the amount of knowledge about the channel state at

the transmitter [we will always assume that the receiver has full channel-state information (CSI)]. If the transmitter has no CSI, then all it can do is transmit with constant power. However, if it had access to CSI, then the transmit power could be controlled as a function of the channel to maximize the capacity. The paper [1] (see also [6]) looked at the problem of maximizing ergodic capacity subject to an average power constraint, and showed that the optimal power-control law was *waterfilling* on the inverse of the channel gain (more power is allocated when the channel is good than when the channel is bad). Another problem that suggests itself is to design power-allocation policies that minimize outage probability on a given fading channel. This problem (among others) was addressed in [4], where it was shown that the best power-allocation scheme was to use no transmit power if the channel is below a threshold, and to use *channel inversion* above the threshold (more power is allocated when the channel is bad than when the channel is good).

The power-allocation policies resulting from maximizing ergodic capacity and from minimizing the probability of outage are very different, and represent two ends of the spectrum. A natural question to ask is whether there is an optimization problem that bridges the gap between these two fundamental problems and their resulting optimal power-control schemes. This is exactly the question tackled in [7]. In [7], an optimal power- and rate-allocation problem is considered, where long-term average capacity is optimized with respect to a deterministic power-allocation policy, subject to a constraint on outage along with the standard average power constraint. This additional outage constraint was motivated by the idea that in an integrated network, non-real-time applications will benefit from maximizing the ergodic capacity, and at the same time, real-time applications (such as voice and video) will benefit from a quality of service (QoS) guarantee on the maximum outage probability. The optimal power allocation for this problem was shown to be a mixture of channel inversion and waterfilling allocation. Extensions of this problem to parallel fading channels with random power-allocation policies (to include discrete fading distributions) have been considered in [8].

In our paper, we generalize these results to a class of fading channels that have a two-time-scale nature. The slow variation in the fading channel is due to distance-based attenuation and shadow fading, and the resulting slow-fading channel gain is known at both transmitter and receiver. The fast-fading gain (resulting from local mobility and multipath fading) is assumed to be known at the receiver but unknown at the transmitter, however, the transmitter does have access to the statistics of the fast fading. In this paper, we restrict our discussion to the case of fast Rayleigh fading where the fast-fading gain is exponentially

Paper approved by M. Chiani, the Editor for Transmission Systems of the IEEE Communications Society. Manuscript received November 18, 2003; revised September 15, 2004. This work was supported by the Australian Research Council (ARC). The Centre for Ultra-Broadband Information Networks (CUBIN) is an affiliated program of the National ICT Australia. This paper was presented in part at the International Symposium on Modeling and Optimization in Mobile, Ad Hoc, and Wireless Networks, Cambridge, U.K., March 2004.

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Digital Object Identifier 10.1109/TCOMM.2005.844935

distributed. Under long codeword and sufficiently long transmission-time assumptions (as in [7]), we define a block-ergodic capacity (BEC) for each code block by averaging over the fast fading. Then we consider an optimum power- and rate-allocation problem, where the long-term average of this BEC (averaged over the slow-fading component) is maximized subject to an average power constraint and an outage constraint on the BEC. The transmitted power is now a function of the slow-fading gain and the statistics of the fast-fading component. We also provide results for the related ergodic capacity and outage capacity problems.

We show that the optimum power-allocation policy is a combination of a “soft waterfilling” policy and channel inversion. We show that (similar to [7]) the resulting long-term average capacity achieves a nice compromise between the corresponding ergodic capacity and outage capacity. We also provide a suboptimal solution to the problem which has a simple power-allocation policy, but achieves near-optimal performance, verified through simulation studies. While the proof techniques are similar to those of [7], there are a number of unique contributions made in this paper.

- Derivation of a “soft waterfilling” power-allocation policy that maximizes ergodic capacity for a two-time-scale fading channel with fast Rayleigh fading.
- Derivation of an optimum power-allocation policy for the problem of maximizing ergodic capacity for a two-time-scale fading channel with fast Rayleigh fading, subject to an average power constraint and an outage constraint.
- Derivation of a simple suboptimal policy that results in near-optimal performance (demonstrated through simulation studies).
- Proofs of a number of inequalities involving the “exponential integral” not reported anywhere else, that are used to derive some intermediate results of this paper.

Before proceeding, we will discuss some related literature dealing with two-time-scale fading, and also the connection of our work to the problem of designing transmission-control schemes based on partial CSI.

The separation of the fading channel into the product of two processes evolving independently at different time scales is standard in wireless channel modeling [9]. The design of power-allocation policies based on knowledge of the slow fading has been considered by many authors (see, for example, [10]–[14] and references therein). Most of these papers deal with cellular systems and consider problems related to meeting average signal-to-interference ratio (SIR) or bit-error rate (BER) constraints, or constraints on the outage probability (which, in this case, means the probability that the SIR drops below a threshold, or equivalently, the BER rises above a threshold).

The problem of two-time-scale fading where the transmitter has access to the slow-fading component only is an example of partial CSI at the transmitter. Many authors have considered problems where the transmitter has partial CSI (see, for example, [15]–[20]). This partial CSI could take the form of a noisy estimate of the channel, a quantized version of the

channel, or channel mean and covariance information in a multiple-antenna setting, to name a few. In our case, the slow-fading gain is the short-term average power of the fast-fading process, and the transmitter thus has knowledge of the local statistics of the fast fading, rather than the fast-fading gain itself. In our case, the local statistics are themselves modeled as a random process, and we look at capacities averaged over this variation.

The rest of the paper is organized as follows. Section II describes the two-time-scale fading-channel model that we use in this paper. Various capacity notions for this fading-channel model are also introduced. We also provide the problem statements for the various different optimal power-allocation tasks we are interested in. Section III presents solutions to these optimization problems. Section IV presents some simulation studies to compare the various capacity results achieved by these power-allocation algorithms. Section V presents some concluding remarks. In order to maintain readability, proofs of our results are relegated to the Appendices or excluded by alluding to similar analyses in [7].

II. CHANNEL MODEL AND VARIOUS CAPACITY NOTIONS

In this paper, we work with a two-time-scale version of the block flat-fading additive white Gaussian noise (BF-AWGN) channel [3] where the channel gain h between the transmitter and the receiver is expressed as $h = gf$, $h > 0$. The gain $g (> 0)$ represents the slow variation of the wireless radio channel, and remains constant over a block of N symbols (one block is spanned by one codeword and N is assumed to be large), but varies from block to block. The gain $f (> 0)$ represents the fast variation of the channel and varies at a much faster rate than the rate of codeword blocks, which implies that the block length is much bigger than the coherence time of the channel. We assume that g has a continuous distribution function $Q(g)$, is ergodic over the time scale of the application concerned, and is independent of the fast-fading process f . The fast-fading process f is assumed to be exponentially distributed (assuming Rayleigh fading) with $E(f) = 1$. It is reasonable to assume that both transmitters and receivers have perfect knowledge of the slow-fading gain g , but the transmitter does not have the knowledge of f . The transmitter power is assumed to be a deterministic function of g , and is denoted by $p(g)$. Thus, it makes sense to define the following conditional maximum achievable rate $R_b(gp(g))$ over each block, termed the BEC

$$R_b(gp(g)) = E \left[\frac{1}{2} \log_2 \left(1 + \frac{gfp(g)}{\sigma^2} \right) \middle| g \right] \quad (1)$$

where σ^2 is the variance of the background white noise, and the conditional expectation E denotes the expectation over the distribution of the fast-fading process f , given g . In the following, “log” will denote natural logarithm.

Lemma 2.1: The BEC is given as

$$R_b(gp(g)) = b_0 e^x E_1(x) \quad x = \frac{\sigma^2}{gp(g)} \quad (2)$$

where $b_0 = 1/(2 \log 2)$ and

$$E_1(x) = \int_x^\infty \frac{e^{-t}}{t} dt.$$

Remark 1: Notice that the BEC $R_b(gp(g))$ is an increasing function of $p(g)$, as one would intuitively expect.

The various capacity notions that are used in the paper are presented below.

- *Ergodic capacity:* The ergodic capacity is defined as the expected value of the BEC, the expectation being taken over the probability density function (pdf) $q(g)$ of the slow-fading gain g . In other words, ergodic capacity

$$R_e = E[R_b(gp(g))] = b_0 E[e^x E_1(x)] \quad x = \frac{\sigma^2}{gp(g)}. \quad (3)$$

- *Outage capacity:* The outage capacity is defined as the maximum achievable BEC over any slow-fading block, denoted by $C_e^b(P_{av})$, where

$$C_e^b(P_{av}) = \max_{p(g)} r \quad (4)$$

subject to the constraints $P(R_b(gp(g)) < r) \leq \epsilon$, $E[p(g)] \leq P_{av}$, and $p(g) \geq 0$.

- The *outage probability* for a given BEC r_0 is defined as

$$O(p(g)) = P(R_b(gp(g)) < r_0). \quad (5)$$

In what follows, we shall solve the following optimization problems.

- P1) Maximize R_e with respect to $p(g)$ subject to $E[p(g)] \leq P_{av}$, $p(g) \geq 0$.
- P2) Maximize R_e with respect to $p(g)$ subject to $E[p(g)] \leq P_{av}$, $p(g) \geq 0$, and $P(R_b(gp(g)) < r_0) \leq \epsilon$ for a given basic BEC r_0 .
- P3) Maximize r with respect to $p(g)$ subject to $P(R_b(gp(g)) < r) \leq \epsilon$, $E[p(g)] \leq P_{av}$, $p(g) \geq 0$, or alternatively,
- P3a) Minimize the outage probability $P(R_b(gp(g)) < r_0)$ with respect to $p(g)$ given a BEC r_0 , subject to $E[p(g)] \leq P_{av}$, $p(g) \geq 0$.

III. OPTIMUM RATE AND POWER ALLOCATION

In this section, we present the optimum power- and rate-allocation solutions to the problems stated above. We first provide the reader with a ‘‘roadmap’’ of the following Sections III-A, III-B, and III-D that are individually devoted to solving Problems **P1**, **P2**, and **P3**, respectively. A suboptimal solution to Problem **P2** is also provided in Section III-C. It should be noted that the solution to **P1** is needed in order to solve Problem **P2**, therefore, the first subsection presents the solution to Problem **P1**, followed by Section III-B that presents the solution to Problem **P2**. Problem **P3** (or **P3a**) can be solved independently of **P1** or **P2**. However, in the single-channel case, the solution to **P3** is rather straightforward, and is given briefly in the last subsection. Recall that **P1** finds the optimal power-control solution to maximize ergodic capacity subject to an average power constraint, **P2** finds the optimal power-control solution to maximize ergodic capacity subject to an average power and a service outage constraint, and **P3** finds the optimal power-control solution to maximize outage capacity (given a specified level of service outage) subject to an average power constraint. Section III-C presents a suboptimal solution to Problem **P2** that is easier to implement than the optimal solution in Sec-

tion III-B. Later, all these different capacity results achieved by the power-control solutions presented in the various subsections are compared via extensive numerical and simulation studies. Below, each individual subsection carries an introductory paragraph on an overview of the rationale behind the solution methodologies used and the relevance of the technical results derived in that subsection.

A. Problem **P1**

In this section, we provide a solution to the optimal power-control problem **P1** using Lagrange-multiplier-based constraint-optimization techniques. Recall that in the absence of fast fading, it is obvious that the solution to **P1** is the well-known *waterfilling solution* [21]. In order to apply the usual Lagrangian-based constrained-optimization theory, we first show (in *Lemma 3.1*) that the objective function in Problem **P1** is a concave function of $p(g)$ (*note*: it is easy to see that the constraint functions are linear in $p(g)$). It is then verified that the optimal power-control solution to **P1** is obtained as a unique fixed-point solution (as a deterministic function of the slow-fading gain g , when g is above a certain threshold) to an implicit equation, and can be computed numerically using an iterative algorithm with guaranteed convergence. It is also seen that when the slow-fading gain falls below this threshold (which is determined by the average power available), the optimal strategy is to turn off transmission. Some key properties of this power-control policy are then presented, as they will be useful in subsequent mathematical analyses. Noticing the similarity between this solution and the well-known *waterfilling* power-control policy for the slow-fading channel case, we call the optimal power-control solution to Problem **P1** *soft waterfilling*. We conclude this subsection with a few comments on the differences between these two policies.

As discussed above, we need the following result in order to proceed (see Appendix III-A for a proof).

Lemma 3.1: The BEC $R_b(g(p(g)))$ is a concave function of $p(g)$.

Remark 2: One can, of course, deduce the above result simply by appealing to the fact that $\log(1 + (gf p(g)/\sigma^2))$ is a concave function of $p(g)$, and the expectation of this function taken over the distribution of f (under the Rayleigh fading assumption) will retain the concavity, a result that follows by the *Dominated Convergence Theorem*.

Lemma 3.1 indicates clearly that R_e is also a concave function of $p(g)$, again by appealing to the Dominated Convergence Theorem, under some mild conditions satisfied by the distribution of the slow fading gain g .

Recalling that the constraints of **P1** are linear in $p(g)$, one can now apply the usual Lagrangian-based constrained-optimization method to form the Lagrangian

$$L(p(g), g, \lambda_0) = \left[R_b(gp(g)) - \lambda_0 \frac{b_0}{\sigma^2} p(g) \right] q(g)$$

where $q(g)$ is the pdf of g and $\lambda_0 > 0$.

Lemma 3.2: The fixed-point equation

$$p^*(g, \lambda_0) = \frac{\sigma^2}{\lambda_0} (1 - ye^y E_1(y)), \quad y = \frac{\sigma^2}{gp^*(g, \lambda_0)} \quad (6)$$

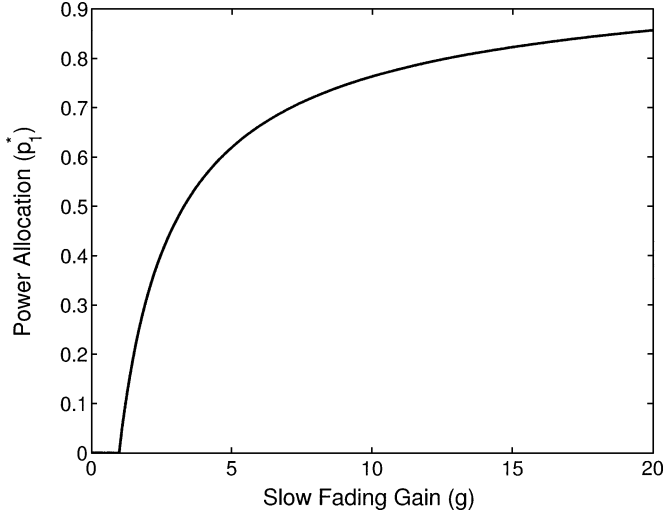


Fig. 1. Plot of soft waterfilling allocation versus the slow-fading gain.

has a unique *positive* fixed-point solution for every $g > \lambda_0$.

Remark 3: The unique fixed-point solution to (6) (for a given value of g) can be computed numerically via an iterative algorithm that is convergent. The details of this algorithm should be obvious from the proof of *Lemma 3.2* (see Appendix II), and are excluded.

Definition 3.1: Suppose for any given $g > \lambda_0$, $p_1^*(g, \lambda_0) > 0$ is the unique fixed point of (6). Define $p_1^*(g, \lambda_0) = 0$ for $0 \leq g \leq \lambda_0$.

It can be easily verified that $p_1^*(g, \lambda_0)$ satisfies the Kuhn–Tucker conditions for optimality (see [22, p. 74] for details). Hence, the solution to the optimization problem **P1** is given by $p_1^*(g, \lambda_0)$, as defined above. The multiplier λ_0 is evaluated by solving $E[p_1^*(g, \lambda_0)] = P_{\text{av}}$. Some key properties of $p_1^*(g, \lambda_0)$ are summarized below.

Lemma 3.3: The optimal power-allocation solution to Problem **P1** [$p_1^*(g, \lambda_0)$] satisfies the following properties: 1) $p_1^*(g, \lambda_0)$ is a strictly increasing function of g for $g > \lambda_0$; and 2) $p_1^*(g, \lambda_0) \rightarrow 0$ as $g \rightarrow \lambda_0$ from above, and $p_1^*(g, \lambda_0) \rightarrow (\sigma^2/\lambda_0)$ as $g \rightarrow \infty$.

Now we discuss some properties of this optimal solution $p_1^*(g, \lambda_0)$, in comparison with the waterfilling solution one would obtain in the absence of fast fading. Recall that the waterfilling solution is given by $p_w^*(g, \lambda_w) = \sigma^2((1/\lambda_w) - (1/g))^+$, where λ_w is the Lagrange multiplier that solves $E[p_w^*(g, \lambda_w)] = P_{\text{av}}$ and $(x)^+ = \max(0, x)$. Both $p_1^*(g, \lambda_0)$ and $p_w^*(g, \lambda_w)$ are zero for values of g below a threshold. In addition, they are both increasing functions of g above this threshold, and approach a limit as $g \rightarrow \infty$. We have observed that $p_1^*(g, \lambda_0)$ tends to increase to the limit more gradually than the waterfilling solution, so, for want of a better name, we call the solution to Problem **P1** “*soft*” waterfilling.

Fig. 1 shows a plot of $p_1^*(g, \lambda)$ versus the slow fading gain g for $\lambda = 1$ with $\sigma^2 = 1$.

B. Problem P2

In this section, we derive the optimum power-allocation policy for Problem **P2**. This problem, as correctly pointed

out in [7], is difficult because of the probabilistic constraint involving the outage probability. One way to circumvent this problem is to consider a probabilistic power-allocation policy, as in [4]. In this paper, however, we focus on a deterministic power-allocation policy. As in [7], we proceed by first deriving a solution to a similar problem (Problem **P2a**), but with a deterministic constraint on the BEC $R_b(gp(g))$. This constraint requires that the BEC is greater than the minimum required rate r_0 over an arbitrary set of the slow-fading gain g . This problem can be solved using standard Lagrangian-multiplier-based techniques for convex optimization, as the BEC is a concave function of $p(g)$. The solution to Problem **P2a** not only provides us with a feasibility condition for Problem **P2**, but also provides us with important insight into how to obtain an optimal solution to **P2**. It is seen that the optimal power-allocation policy for **P2** is analogous to the solution to **P2a**, except that this solution guarantees that $R_b(gp(g)) \geq r_0$ over a particular set of the slow-fading gain determined by the maximum service outage probability ϵ . *Theorem 1* summarizes the solution to **P2a**, while *Theorem 2* provides an intermediate result. These two theorems (see Appendices for proofs) together lead to the main result of this paper, namely, the optimal solution to Problem **P2**, following a similar analysis in [7]. This result is stated in *Theorem 3*, which shows that the optimal power-allocation policy for **P2** is a combination of channel inversion and the optimal power-allocation solution to Problem **P1**, namely, the soft waterfilling policy. The rest of the subsection provides the required mathematical analysis. Proofs are either relegated to the Appendices or excluded by alluding to a similar analysis in [7].

In order to proceed, we need the definition of a *service set* and an *outage set*, as in [7].

Definition 3.2: Given a basic rate r_0 and a power policy $p(g)$, the *service set* is defined as

$$\mathcal{H}_s(p(g)) = \{g | R_b(gp(g)) \geq r_0\}.$$

Correspondingly, the *outage set* is defined as

$$\mathcal{H}_o(p(g)) = \{g | R_b(gp(g)) < r_0\}.$$

It will be seen that our solution to Problem **P2** will result in a power-allocation policy having a particular form of a service set (mirroring a similar result in [7]). In order to derive this result, we first solve the following problem, where it is required that the service set contains an arbitrary set \mathcal{H}_r .

Problem P2a: Find $R^*(\mathcal{H}_r) = \max_{p(g)} R_e$ subject to

$$\begin{aligned} E(p(g)) &\leq P_{\text{av}} \\ p(g) &\geq 0 \\ R_b(gp(g)) &\geq r_0 \quad \forall g \in \mathcal{H}_r. \end{aligned}$$

It is obvious (by noting that $e^x E_1(x)$ is a monotonically decreasing function of x) that a feasible solution $p(g)$ to Problem **P2a** has to satisfy

$$p(g) \geq \frac{\sigma^2}{gx_0} \quad g \in \mathcal{H}_r$$

where $e^{x_0} E_1(x_0) = r_0/b_0$. This implies that for Problem **P2a** to be feasible, a necessary condition is $P_{\text{av}} \geq P_{\text{min}}(r_0, \mathcal{H}_r)$ where

$$P_{\text{min}}(r_0, \mathcal{H}_r) = \int_{\mathcal{H}_r} \frac{\sigma^2}{gx_0} q(g) dg. \quad (7)$$

Let us denote an optimal solution to Problem **P2a** by $p^*(g, \mathcal{H}_r)$. Following the usual Lagrange-multiplier-based convex-optimization theory, one obtains the following solution to Problem **P2a** (see Appendix IV for a proof).

Theorem 1: Suppose $P_{\text{av}} = P_{\text{min}}(r_0, \mathcal{H}_r)$. Then the optimal power-allocation solution to Problem **P2a** is given by

$$p^*(g, \mathcal{H}_r) = \begin{cases} \frac{\sigma^2}{gx_0}, & g \in \mathcal{H}_r \\ 0, & \text{otherwise.} \end{cases} \quad (8)$$

If $P_{\text{av}} > P_{\text{min}}(r_0, \mathcal{H}_r)$, then the optimal power-allocation solution to Problem **P2a** is given by

$$p^*(g, \mathcal{H}_r) = \begin{cases} \frac{\sigma^2}{gx_0}, & g \in \mathcal{H}_r \cap \{g \leq \bar{g}\} \\ p_1^*(g, \lambda_1), & \text{otherwise} \end{cases} \quad (9)$$

where $\bar{g} = \lambda_1/(x_0(1 - x_0 r_0/b_0))$, and λ_1 is the solution to $E[p^*(g, \mathcal{H}_r)] = P_{\text{av}}$.

The interpretation of the result above is that the optimal power allocation is a combination of channel inversion over the slow-fading component and the optimal power allocation for obtaining maximum average block-ergodic rate given by the solution to Problem **P1**, namely, the soft waterfilling solution. Within the service set, the allocated power is no less than that achieved by inverting the slow-fading component to maintain the basic block-ergodic rate r_0 . For slow-fading gains larger than a threshold (\bar{g}), the allocated power is larger, given by the soft waterfilling solution to achieve the maximum average block-ergodic rate.

As in [7], we now define the set of good channels and bad channels according to the value of the slow-varying channel gain g . Defining $Q(g_\epsilon) = \epsilon$, the set of good channels is defined as $\mathcal{H}_\epsilon = \{g : g \geq g_\epsilon\}$, and the set of bad channels is defined as $\bar{\mathcal{H}}_\epsilon = \{g : g < g_\epsilon\}$. We will show that the solution to Problem **P2a** with $\mathcal{H}_r = \mathcal{H}_\epsilon$, i.e., $p^*(g, \mathcal{H}_\epsilon)$, is an optimum solution to Problem **P2**.

Define the partial ordering of two sets, as in [7], that is, $\mathcal{H}_1 \prec \mathcal{H}_2$ if $g_1 < g_2, \forall g_1 \in \mathcal{H}_1$ and $g_2 \in \mathcal{H}_2$. Then one can prove the following theorem.

Theorem 2: Problem **P2** has an optimum solution $p^*(g) = p^*(g, \mathcal{H}_\epsilon)$ with the outage set $\mathcal{H}_o(p^*(g))$ and service set $\mathcal{H}_s(p^*(g))$, such that $\mathcal{H}_o(p^*(g)) \prec \mathcal{H}_s(p^*(g))$.

Proof: The proof is based on a complex two-stage construction process, and is relegated to Appendix V. ■

This theorem basically mirrors the corresponding result in [7, Th. 2] in that it implies that the block-ergodic rate falls below the basic rate r_0 when the slowly varying channel gain falls below a particular threshold.

Now using the fact that $P(\mathcal{H}_s(p^*(g))) \geq 1 - \epsilon$, one can easily show that Problem **P2** has an optimum solution $p^*(g) = p^*(g, \mathcal{H}_\epsilon)$, such that $\mathcal{H}_\epsilon \subseteq \mathcal{H}_s(p^*(g))$. A similar analysis to [7] now leads to the main result of this paper.

Theorem 3: Problem **P2** has a solution if and only if $(P_{\text{av}}, r_0, \epsilon)$ satisfies $P_{\text{av}} \geq P_{\text{min}}(r_0, \mathcal{H}_\epsilon)$. When $P_{\text{av}} = P_{\text{min}}(r_0, \mathcal{H}_\epsilon)$, the optimum power allocation is given by

$$p^*(g) = \begin{cases} \frac{\sigma^2}{gx_0}, & g \in \mathcal{H}_\epsilon \\ 0, & \text{otherwise.} \end{cases} \quad (10)$$

When $P_{\text{av}} > P_{\text{min}}(r_0, \mathcal{H}_\epsilon)$, an optimum power allocation is given by

$$p^*(g) = \begin{cases} \frac{\sigma^2}{gx_0}, & g_\epsilon \leq g \leq \max(g_\epsilon, g^*) \\ p_1^*(g, \lambda^*), & \text{otherwise} \end{cases} \quad (11)$$

where $\bar{g}^* = \lambda^*/x_0(1 - x_0 r_0/b_0)$, and λ^* is the solution to $E[p^*(g, \mathcal{H}_\epsilon)] = P_{\text{av}}$.

Note that computing $p^*(g)$ involves computing $p_1^*(g, \lambda^*)$, which has to be computed numerically (for a given value of g) as the unique fixed-point solution to an implicit equation, as given by (6). In order to avoid this computational complexity, one may be motivated to look at a suboptimal solution to Problem **P2**. This is the topic of the following subsection.

C. A Suboptimal Solution to Problem **P2**

In this section, we present a simple suboptimal solution to Problem **P2** by noticing that

$$E \left[E \left[\log \left(1 + \frac{gfp(g)}{\sigma^2} \right) \middle| g \right] \right] \leq E \left[\log \left(1 + \frac{gp(g)}{\sigma^2} \right) \right]$$

by Jensen's inequality, remembering that $E[f] = 1$. Now one can pose the following optimization task.

Problem P2b: Maximize $E[\log(1 + (gp(g)/\sigma^2))]$ with respect to $p(g)$, subject to $E[p(g)] \leq P_{\text{av}}$ and $p(g) \geq 0$ and $P(R_b(gp(g)) < r_0) \leq \epsilon$.

Notice that this problem can be treated in a similar way to that of [7] by observing that the constraint $P(R_b(gp(g)) < r_0) \leq \epsilon$ can be rewritten as $P(p(g)) < (\sigma^2/gx_0) \leq \epsilon$.

An optimal solution to Problem **P2b** can now be obtained in exactly the same fashion as in [7] (hence, we do not provide the analysis), and we call this solution $p_{\text{sub}}(g)$. The next theorem states this result.

Theorem 4: A suboptimal solution to Problem **P2**, given as an optimum solution to **P2b**, is given by [if $P_{\text{av}} = P_{\text{min}}(r_0, \mathcal{H}_\epsilon)$]

$$p_{\text{sub}}(g) = \begin{cases} \frac{\sigma^2(2^{2r_1}-1)}{g}, & g \in \mathcal{H}_\epsilon \\ 0, & \text{otherwise.} \end{cases} \quad (12)$$

Otherwise, if $P_{\text{av}} > P_{\text{min}}(r_0, \mathcal{H}_\epsilon)$

$$p_{\text{sub}}(g) = \begin{cases} \frac{\sigma^2(2^{2r_1}-1)}{g}, & g_\epsilon \leq g \leq \max(g_\epsilon, g_{\text{sub}} * 2^{2r_1}) \\ \sigma^2 \left(\frac{1}{g_{\text{sub}}} - \frac{1}{g} \right)^+, & \text{otherwise} \end{cases} \quad (13)$$

where $r_1 = b_0 \log(1 + (1/x_0))$ and g_{sub} is the solution to $E[p_{\text{sub}}(g)] = P_{\text{av}}$.

Later, we will compare the capacity results achieved by the optimal and suboptimal solutions to **P2** via simulation studies.

D. Problem **P3**

In this section, we provide the optimal solution to Problems **P3a** and **P3**. Note that the proof is intuitively straightforward in the case of a single communication channel, with the

two-time-scale fading properties satisfying the assumptions in Section II, and the power-allocation policy being deterministic. Because of this, we only provide a brief sketch of a proof of *Theorem 5*. In addition, by noting that the service outage probability can also be written as $P(gp(g) < (\sigma^2/x_0))$ (where $e^{x_0}E_1(x_0) = r_0$), it can be easily seen that the solution to Problem **P3** can be obtained as a special case of the outage-capacity maximization problem for multiple parallel channels considered in [4], where the number of channels is one. The corresponding optimal power-allocation problem for outage capacity maximization, in the case of multiple parallel communication channels with two-time-scale fading properties (where we can also allow the slow-fading gain to be a discrete random process and the power-allocation policy to be probabilistic), is complex and is beyond the scope of this paper. This problem will be treated rigorously in a subsequent paper.

Theorem 5: The outage probability is minimized by the optimal power-allocation solution to Problem **P3a**

$$p_{\text{out}}(g) = \begin{cases} \frac{\sigma^2}{gx_0}, & g > g_{\text{out}} \\ 0, & \text{otherwise} \end{cases} \quad (14)$$

where $\int_{g_{\text{out}}}^{\infty} (\sigma^2/gx_0)q(g)dg = P_{\text{av}}$. Similarly, the outage capacity is given by \bar{r} , such that $\bar{r} = b_0 \exp(\bar{x})E_1(\bar{x})$ and \bar{x} is given by $\int_{g_{\text{out}}}^{\infty} (\sigma^2/g\bar{x})q(g)dg = P_{\text{av}}$.

Proof: It is easy to show that the BEC within the service set should be exactly r_0 , since if it is greater than r_0 for some subset of the service set, the service set could be expanded by redistributing power [since the BEC is an increasing function of $p(g)$], such that the average power constraint is satisfied, thus reducing outage probability. Similarly, the optimal power-allocation solution $p_{\text{out}}(g)$ must also satisfy the average power constraint with equality, otherwise additional power could be invested in enlarging the service set. The particular form of the solution $p_{\text{out}}(g)$, as given above by *Theorem 5*, thus follows easily. ■

IV. SIMULATION RESULTS

In this section, we present some simulation studies conducted with a single-user channel. The slow-fading gain g is distributed with a lognormal distribution (standard for shadow fading), such that $\log g$ is distributed with mean 0 and variance $\sigma_g = 10^{0.3}$. The fast-fading gain f is exponentially distributed with mean 1 (Rayleigh fading). The minimum basic rate is taken to be $r_0 = 1$ b/s/Hz. The maximum outage probability is constrained to be $\epsilon = 0.05$. (P_{av}/σ^2) is varied between 11 and 14 dB.

Fig. 2 shows how the various capacity measures increase with increasing signal-to-noise ratio (SNR). The SNR is computed as $10 \log_{10}(P_{\text{av}}/\sigma^2)$. As it can be seen, capacity with outage achieves a compromise between ergodic capacity and outage capacity, the gap between ergodic capacity and capacity with outage decreasing with increasing average transmission power. What is interesting is to observe that the suboptimal power-allocation solution $p_{\text{sub}}(g)$ results in nearly optimal performance. Thus, one can use the easily computable solution $p_{\text{sub}}(g)$ (as opposed to $p^*(g)$) in order to achieve good capacity performance with the same outage constraint.

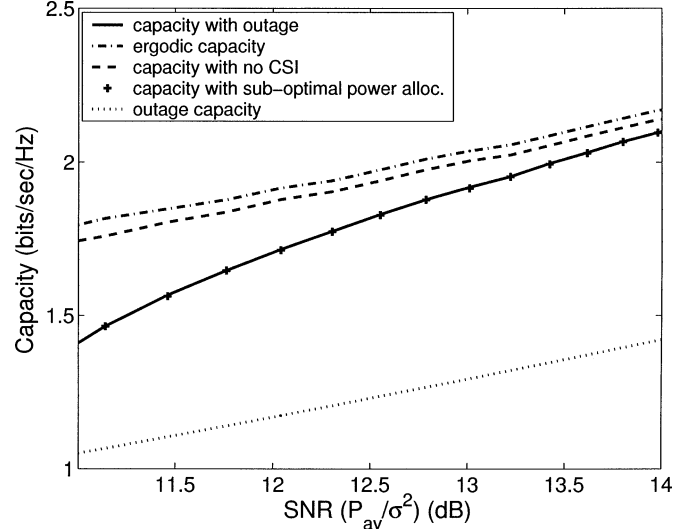


Fig. 2. Various capacity plots against P_{av}/σ^2 .

V. CONCLUSION

In this paper, we have studied the power-control problem for wireless channels from an information-theoretic perspective. Integral to the developments was the use of a two-time-scale fading model which incorporated slow fading due to distance and shadowing effects, and fast fading due to multipath propagation. We assumed that the transmitter had knowledge of the slow-fading gain and the statistic of the fast fading, but did not know the instantaneous value of the fast-fading component. Power-control schemes were thus designed to exploit knowledge of the slow fading. We considered various notions of *capacity*, including ergodic capacity, outage capacity, and capacity subject to a service outage constraint, developing optimal power-allocation strategies in each case.

Extensions of this work are currently being considered for parallel channels, as well as to the case of multiaccess channels with applications to opportunistic scheduling.

APPENDIX I

PROOF OF LEMMA 3.1

In this proof, to keep notations simple, we use R_b instead of $R_b(gp(g))$. It is easy to work out that [with $x = (\sigma^2/gp(g))$]

$$\frac{d^2 R_b}{dp^2(g)} = -\frac{b_0}{p^2(g)} \left[(1 - xe^x E_1(x)) + p(g) \{ (x+1)e^x E_1(x) - 1 \} \right].$$

Using the fact that $p(g) \geq 0$, and the two standard inequalities $e^x E_1(x) \leq (1/x)$, $e^x E_1(x) > (1/(x+1))$ for $x > 0$, it follows that $(d^2 R_b/dp^2(g)) < 0$. Thus $R_b(gp(g))$ is a concave function of $p(g)$.

APPENDIX II

PROOF OF LEMMA 3.2

Notice that (6) can be written as

$$\frac{\lambda_0}{g} = y [1 - ye^y E_1(y)].$$

Define $q_1(y) = y[1 - ye^y E_1(y)]$. First, one can easily check that for $y = 0$, $q_1(y) = 0$. Notice that

$$\frac{dq}{dy} = 1 + y - ye^y E_1(y)(y + 2).$$

Using the continued fraction expansion of $e^y E_1(y)$ (see [23]), one can write for $y > 0$

$$ye^y E_1(y) = \frac{1}{1 + \frac{1}{y + \frac{1}{1 + \dots}}} < \frac{1}{1 + \frac{1}{y+1}} = \frac{y+1}{y+2}$$

which implies that $(dq/dy) > 0$, for $y > 0$. Therefore, $q_1(y) = y[1 - ye^y E_1(y)]$ is a strictly monotonically increasing function for $y > 0$. Also, it is straightforward to show that $q_1(y) < 1$ for $y > 0$ and $q_1(y) \rightarrow 1$ as $y \rightarrow \infty$. Therefore, for every given $g > \lambda_0$, (6) has a unique fixed-point solution $p^*(g, \lambda_0) > 0$.

APPENDIX III PROOF OF LEMMA 3.3

1) We need the following proposition involving the exponential integral function $E_1(x)$.

Proposition 1: For $x > 0$

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)}$$

is an increasing function of x . As $x \rightarrow \infty$

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} \rightarrow 1.$$

It also follows that

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} < 1$$

for $0 < x < \infty$.

Proof: In this proof, we assume $x > 0$. Let us denote

$$V(x) = \frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)}.$$

By evaluating $dV(x)/dx$ one gets that a necessary condition for $dV(x)/dx > 0$ is

$$\left[e^x E_1(x) - \frac{3x+1}{2x^2} \right]^2 < \frac{x^2 + 6x + 1}{4x^4}$$

which implies [noting that $e^x E_1(x) < (1/x)$]

$$\frac{3x+1}{2x^2} > e^x E_1(x) > \frac{3x+1 - \sqrt{x^2 + 6x + 1}}{2x^2}.$$

The inequality of the left-hand side (LHS) is trivial to prove, since $e^x E_1(x) < (1/x)$. The right-hand side (RHS) inequality is rather tricky to prove. One proof based on a continued fraction expansion of $e^x E_1(x)$ [24] is as follows:

$$e^x E_1(x) = \frac{1}{x + \frac{1}{1 + \frac{1}{x + \frac{1}{1 + \frac{1}{x + \frac{1}{1 + \frac{1}{x + \frac{1}{1 + \frac{1}{x + \frac{1}{1 + \dots}}}}}}}}}}}. \quad (15)$$

It is then straightforward to show (working upwards from the lowest level shown) that

$$\begin{aligned} e^x E_1(x) &> \frac{x^3 + 15x^2 + 58x + 50}{x^4 + 16x^3 + 72x^2 + 96x + 24} \\ &> \frac{(3x+1) - \sqrt{x^2 + 6x + 1}}{2x^2} \end{aligned}$$

where the first inequality follows from the continued fraction expansion, and the second inequality follows after some tedious algebraic manipulation. This implies that

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)}$$

is an increasing function of x for $x > 0$.

Notice also trivially that

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} > 0$$

for $x > 0$. Using the asymptotic expansion for $e^x E_1(x)$ [23], one can show that

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} \rightarrow 1$$

as $x \rightarrow \infty$. This also implies that

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} < 1$$

for $x > 0$. ■

From Lemma 3.2, $p_1^*(g, \lambda_0)$ satisfies the following equation for $g > \lambda_0$:

$$p_1^*(g, \lambda_0) = \frac{\sigma^2}{\lambda_0} (1 - xe^x E_1(x)) \quad x = \frac{\sigma^2}{gp_1^*(g, \lambda_0)}. \quad (16)$$

This implies that for $g > \lambda_0$

$$\begin{aligned} \frac{dp_1^*(g, \lambda_0)}{dg} &= \frac{\sigma^2}{\lambda_0} [1 - (x+1)e^x E_1(x)] \frac{dx}{dg} \\ &= \frac{x^2}{\lambda_0} [(x+1)e^x E_1(x) - 1] \\ &\quad \times \left(p_1^*(g, \lambda_0) + g \frac{dp_1^*(g, \lambda_0)}{dg} \right). \end{aligned}$$

Rearranging, one can write

$$\begin{aligned} \frac{dp_1^*(g, \lambda_0)}{dg} &\left[1 - \frac{x^2}{\lambda_0} \{(x+1)e^x E_1(x) - 1\} g \right] \\ &= \frac{x^2}{\lambda_0} p_1^*(g, \lambda_0) \{(x+1)e^x E_1(x) - 1\}. \end{aligned}$$

Since $e^x E_1(x) > (1/(x+1))$, it is easy to see that the RHS of the above equation is positive. We now show that

$$\left[1 - \frac{x^2}{\lambda_0} \{(x+1)e^x E_1(x) - 1\} g \right] > 0.$$

By rearranging (16), one can easily show that the LHS of the previous inequality is identical to

$$1 - \frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)}.$$

It can be shown that

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} < 1$$

for $x > 0$ (see *Proposition 1*). Thus, $(dp_1^*(g, \lambda_0)/dg) > 0$ for $g > \lambda_0$ and 1) is proved.

2) As g increases, x decreases [from part 1)], hence $x e^x E_1(x) \rightarrow 0$ as $g \rightarrow \infty$. Therefore, $p_1^*(g, \lambda_0) \rightarrow (\sigma^2/\lambda_0)$.

Similarly, as $g \rightarrow \lambda_0$ from above, $p_1^*(g, \lambda_0)$ decreases strictly monotonically. As $p_1^*(g, \lambda_0) \geq 0$, $p_1^*(g, \lambda_0) \rightarrow 0$ as $g \rightarrow \lambda_0$.

APPENDIX IV PROOF OF THEOREM 1

When $P_{\text{av}} = P_{\min}(r_0, \mathcal{H}_r)$, solution (8) is obvious. When $P_{\text{av}} > P_{\min}(r_0, \mathcal{H}_r)$, the problem can be translated to the following problem.

Maximize $E[R_b(gp(g))]$ with respect to $p(g)$ subject to $E[p(g)] = P_{\text{av}}$, $p(g) \geq 0$ and $p(g) \geq (\sigma^2/gx_0)$, $g \in \mathcal{H}_r$.

Notice that from *Lemma 3.1*, $R_b(gp(g))$ is a concave function of $p(g)$. Therefore, Problem **P2a** has a concave objective function and linear or concave constraints. Then $p^*(g, \mathcal{H}_r)$ is the optimal solution iff it satisfies the Kuhn–Tucker conditions [25].

Using a Lagrange multiplier $b_0 \lambda_1 > 0$, we define the Lagrangian

$$L(p(g), g, \lambda_1) = [R_b(gp(g)) - \lambda_1 b_0 p(g)] q(g).$$

Define $\tilde{\mathcal{H}}_g$ to be the union of the sets $\{g \in \mathcal{H}_r | p^*(g, \mathcal{H}_r) > (\sigma^2/gx_0)\}$ and $\{g \notin \mathcal{H}_r | p^*(g, \mathcal{H}_r) > 0\}$. It is then straightforward to show that when $g \in \tilde{\mathcal{H}}_g$

$$\left. \frac{dL(p(g), g, \lambda_1)}{dp(g)} \right|_{p(g)=p^*(g, \mathcal{H}_r)} = 0$$

and

$$\left. \frac{dL(p(g), g, \lambda_1)}{dp(g)} \right|_{p(g)=p^*(g, \mathcal{H}_r)} \leq 0$$

when $g \notin \tilde{\mathcal{H}}_g$. Since $p^*(g, \mathcal{H}_r)$ is unique and positive $\forall g > \lambda_1$ (*Lemma 3.2*), and $p^*(g, \mathcal{H}_r) = 0$ for $g \leq \lambda_0$, the Kuhn–Tucker conditions are satisfied. Therefore, $p^*(g, \mathcal{H}_r)$ is the optimum solution to Problem **P2a**.

APPENDIX V PROOF OF THEOREM 2

The proof of this theorem is based on a two-stage construction process provided in [7]. To avoid repetition, we summarize most of the proof. However, there are certain intermediate results needed, which are not straightforward to prove. More details are provided for these results.

Suppose there is an arbitrary feasible power policy for Problem **P2** given by $\hat{p}(g)$ with service set $\hat{\mathcal{H}}_s$ and long-term average rate \hat{R} , with $E[\hat{p}(g)] \leq P_{\text{av}}$ and $P(\hat{\mathcal{H}}_s) \geq 1 - \epsilon$. Then,

another policy $\tilde{p}(g)$ can be constructed by letting $\mathcal{H}_r = \hat{\mathcal{H}}_s$ in Problem **P2a**, such that [for $P_{\text{av}} > P_{\min}(r_0, \hat{\mathcal{H}}_s)$]

$$\tilde{p}(g) = \begin{cases} \frac{\sigma^2}{gx_0}, & g \in \hat{\mathcal{H}}_s \cap \{g \leq \tilde{g}\} \\ p_1^*(g, \tilde{\lambda}), & \text{otherwise} \end{cases} \quad (17)$$

where $\tilde{g} = \tilde{\lambda}/(x_0(1 - x_0 r_0/b_0))$ and $\tilde{\lambda}$ is the solution to $E[\tilde{p}(g)] = P_{\text{av}}$. For $P_{\text{av}} = P_{\min}(r_0, \hat{\mathcal{H}}_s)$

$$\tilde{p}(g) = \begin{cases} \frac{\sigma^2}{gx_0}, & g \in \mathcal{H}_r \\ 0, & \text{otherwise.} \end{cases} \quad (18)$$

$\tilde{p}(g)$ is obviously feasible and achieves a higher long-term average rate than \hat{R} , since $\tilde{p}(g)$ achieves the highest average rate among all policies whose service sets contain $\hat{\mathcal{H}}_s$.

Now define the following power policies over the entire space of the slow-fading gain $g \in (0, \infty)$:

$$\begin{aligned} \tilde{p}_{\text{swf}}(g) &= p_1^*(g, \tilde{\lambda}) \quad 0 < g < \infty \\ \tilde{p}_{\text{res}}(g) &= \left[\frac{\sigma^2}{gx_0} - \tilde{p}_{\text{swf}}(g) \right]^+ \quad 0 < g < \infty. \end{aligned} \quad (19)$$

Also define the set (over which $R_b(g\tilde{p}(g)) \geq r_0$) $\tilde{\mathcal{H}}_{\text{swf}} = \{g : g \geq \tilde{g}\}$ and $\tilde{\mathcal{H}}_{\text{inv}} = \hat{\mathcal{H}}_s \setminus \tilde{\mathcal{H}}_{\text{swf}}$. Notice that $\tilde{p}(g)$ can also be written as

$$\tilde{p}(g) = \tilde{p}_{\text{swf}}(g) + 1_{g \in \tilde{\mathcal{H}}_{\text{inv}}} \tilde{p}_{\text{res}}(g) \quad (20)$$

where 1_x is the indicator function taking value 1 when x is true, and 0 otherwise. Now construct another power-allocation policy

$$p'(g) = \tilde{p}_{\text{swf}}(g) + 1_{g \in \mathcal{H}'_{\text{inv}}} \tilde{p}_{\text{res}}(g) \quad (21)$$

where $\mathcal{H}'_{\text{inv}} = \{g'_b \leq g < \tilde{g}\}$ and

$$\int_{\mathcal{H}'_{\text{inv}}} \tilde{p}_{\text{res}}(g) q(g) dg = \text{int}_{\tilde{\mathcal{H}}_{\text{inv}}} \tilde{p}_{\text{res}}(g) q(g) dg. \quad (22)$$

Denote the average rates achieved by $\tilde{p}(g)$ and $p'(g)$ as \tilde{R} and R' , respectively. Then, we have the following lemma.

Lemma 5.1: The power-allocation policy $p'(g)$ has the following properties:

- 1) $E[p'(g)] = E[\tilde{p}(g)] = P_{\text{av}}$;
- 2) $\mathcal{H}_o(p'(g)) \prec \mathcal{H}_s(p'(g))$;
- 3) $R' \geq \tilde{R}$;
- 4) $P\{\mathcal{H}_s(p'(g))\} \geq P\{\mathcal{H}_s(\tilde{p}(g))\}$.

Proof: Proofs of 1) and 2) are trivial (see [7]) for details. Proofs of 3) and 4) require *Proposition 2* and *Proposition 4*. The proof of *Proposition 2* is algebraically rather complex and is provided in sufficient detail, whereas *Proposition 4* is directly quoted from [7].

Proposition 2: The power-efficiency function given by

$$\eta(g) = \frac{r_0 - R_b(g\tilde{p}_{\text{swf}}(g))}{\tilde{p}_{\text{res}}(g)} \quad 0 < g < \tilde{g} \quad (23)$$

is an increasing function of g .

Proof: Using the fact that $\tilde{p}_{\text{swf}}(g)$ satisfies (6), one can (after some algebraic manipulation) rewrite $1/\eta(g)$ as

$$\beta(x) = c_1 \frac{(x - x_0)(1 - xe^x E_1(x))}{e^{x_0} E_1(x_0) - e^x E_1(x)}$$

where c_1 is a generic constant. For $g < \check{g}$, $x > x_0$ as at $g = \check{g}$, $\tilde{p}_{\text{swf}}(g) = \sigma^2/gx_0$. Now we need to show that $\beta(x)$ is an increasing function of x for $x > x_0$ where $x_0 > 0$. Denoting $\alpha(x) = (x - x_0)(1 - xe^x E_1(x))$ and $\gamma(x) = e^{x_0} E_1(x_0) - e^x E_1(x)$, notice that at $\alpha(x_0) = \gamma(x_0) = 0$. Note also that both $\alpha(x)$ and $\gamma(x)$ are positive for $x > x_0$. It is obvious that $\gamma(x)$ is an increasing function of x for $x > x_0$. The same can be said about $\alpha(x)$, by observing that $(d/dx)\alpha(x) = x[(1/x - e^x E_1(x)) - ((x+1)e^x E_1(x) - 1)] + x_0$, which is positive from *Proposition 1* and the fact $x_0 > 0$.

Now it is straightforward to derive that

$$\frac{\frac{d}{dx}\alpha(x)}{\frac{d}{dx}\gamma(x)} = x \left[1 - \frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} \right] + x_0 \left[\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} \right].$$

Clearly, from *Proposition 1*

$$\frac{\frac{d}{dx}\alpha(x)}{\frac{d}{dx}\gamma(x)} > 0.$$

We also know from *Proposition 1* that

$$\frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)}$$

is an increasing function of x for $x > 0$. Hence, we just need to show that

$$x \left[1 - \frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} \right]$$

is an increasing function of x for $x > 0$, which is the following proposition.

Proposition 3: For $x > 0$

$$x \left[1 - \frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} \right]$$

is an increasing function of x .

Proof: The proof of this is similar to that of *Proposition 1*. Writing

$$W(x) = x \left[1 - \frac{(x+1)e^x E_1(x) - 1}{\frac{1}{x} - e^x E_1(x)} \right]$$

one can (after some algebra) show that $(dW(x)/dx) > 0$ implies

$$(e^x E_1(x))^2 2(x+1) - e^x E_1(x) \left(5 + \frac{4}{x} \right) + \frac{3x+1}{x^2} > 0.$$

This implies

$$\left[e^x E_1(x) - \frac{5x+4}{4x(x+1)} \right]^2 > \frac{x^2+8x+8}{4x(x+1)^2}.$$

Noting again that

$$\frac{5x+4}{4x(x+1)} > \frac{1}{x} > e^x E_1(x)$$

one needs to now show that

$$e^x E_1(x) < \frac{5x+4 - \sqrt{x^2+8x+8}}{4x(x+1)}$$

for $x > 0$. The proof of this is again based on the continued fraction expansion (15), which leads to

$$e^x E_1(x) < \frac{x^3+11x^2+26x+6}{x(x^3+12x^2+36x+24)} < \frac{5x+4 - \sqrt{x^2+8x+8}}{4x(x+1)}.$$

This completes the proof. \blacksquare

Therefore, we have shown that $(d/dx)\alpha(x)/(d/dx)\gamma(x)$ is a positive increasing function of x for $x > x_0$, and the proof is complete. \blacksquare

Proposition 4: For two disjoint sets Ψ_1, Ψ_2 , let $f_1(x)$ be an arbitrary function, such that $f_1(x'') \geq f_1(x') \forall x'' \in \Psi_2, x' \in \Psi_1$. For any nonnegative function $f_2(x)$ satisfying $\int_{\Psi_2} f_2(x) dx = \int_{\Psi_1} f_2(x) dx$, we have

$$\int_{\Psi_2} f_1(x) f_2(x) dx \geq \int_{\Psi_1} f_1(x) f_2(x) dx.$$

Now we continue our proofs of parts 3) and 4) of *Lemma 5.1*.

Proof of 3): Define $\Psi = \mathcal{H}'_{\text{inv}} \cap \mathcal{H}'_{\text{inv}}$ such that $\tilde{\Psi} = \mathcal{H}'_{\text{inv}} \Psi$ and $\Psi' = \mathcal{H}'_{\text{inv}} \Psi$ are two disjoint sets. The average rate of $\tilde{p}(g)$ can be expressed as

$$\tilde{R} = \int_0^\infty R_b(g\tilde{p}_{\text{swf}}(g))q(g)dg + \int_{\mathcal{H}'_{\text{inv}}} (r_0 - R_b(g\tilde{p}_{\text{swf}}(g)))q(g)dg.$$

Using the definition of the power-efficiency function (23), one can write

$$\tilde{R} = R_{\text{swf}} + \int_{\Psi} \eta(g)\tilde{p}_{\text{res}}(g)q(g)dg + \int_{\tilde{\Psi}} \eta(g)\tilde{p}_{\text{res}}(g)q(g)dg$$

where $R_{\text{swf}} = \int_0^\infty R_b(g\tilde{p}_{\text{swf}}(g))q(g)dg$. From a similar expression for R' , we finally get

$$R' - \tilde{R} = \int_{\Psi'} \eta(g)\tilde{p}_{\text{res}}(g)q(g)dg - \int_{\tilde{\Psi}} \eta(g)\tilde{p}_{\text{res}}(g)q(g)dg.$$

Since $\tilde{P}si \prec \Psi'$ (see [7]), by *Proposition 2*, *Proposition 4*, and (22), we get $R' \geq \tilde{R}$.

Proof of 4): By repeating the analysis in [7], it can be shown that $P\{\mathcal{H}_s(p'(g))\} - P\{\mathcal{H}_s(\tilde{p}(g))\} = P(\Psi') - P(\tilde{\Psi})$, which is equal to

$$\int_{\Psi'} \frac{1}{\tilde{p}_{\text{res}}(g)} \tilde{p}_{\text{res}}(g)q(g)dg - \int_{\tilde{\Psi}} \frac{1}{\tilde{p}_{\text{res}}(g)} \tilde{p}_{\text{res}}(g)q(g)dg.$$

Note that $\tilde{p}_{\text{res}}(g) = (\sigma^2/gx_0) - \tilde{p}_{\text{swf}}(g)$ when $g \in \Psi'$ or $g \in \tilde{\Psi}$. Since from *Lemma 3.3*, $\tilde{p}_{\text{swf}}(g)$ is an increasing function of g , it is obvious that $\tilde{p}_{\text{res}}(g)$ is a decreasing function of g , and hence, $1/\tilde{p}_{\text{res}}(g)$ is an increasing function of g . *Proposition 4* now implies that $P\{\mathcal{H}_s(p'(g))\} \geq P\{\mathcal{H}_s(\tilde{p}(g))\}$.

Thus, starting from any arbitrary feasible power allocation $\tilde{p}(g)$, one can construct a better power-allocation policy $p'(g)$, such that $\mathcal{H}_o(p'(g)) < \mathcal{H}_s(p'(g))$. This proves *Theorem 2*. ■

ACKNOWLEDGMENT

The authors would like to thank J. Luo for useful discussions.

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