

# An introduction to Stochastic Optimisation

Emilien Flayac & Alejandro I. Maass

The University of Melbourne



### Broad view on optimisation

Optimisation is about making a choice among a range of possible options according to some quantitative criterion

An optimisation problem is defined by:

- A *feasible* (or *admissible*) set X that represents the range of our possibilities
- A cost (or objective) function  $f : \mathcal{X} \to \mathbb{R}$  that weighs every feasible option

$$\min_{x \in \mathcal{X}} f(x) \tag{P}$$

# Solution $x^*$ is a solution of (P) iff $x^* \in \mathcal{X}$ and $\forall x \in \mathcal{X}$ , $f(x^*) \leq f(x)$

Introduction to Stochastic Optimisation



## Broad view on optimisation: Admissible set

The nature of  $\mathcal X$  depends on the nature of the decision:

- Discrete vs. Continuous
- Finite vs. Infinite dimensional
- Static vs. Dynamic
- Defined by a finite set of inequality constraints:

$$\mathcal{X} = \{ x \in E : c(x) \le 0 \}$$

with  $c(x) = (c_1(x), ..., c_m(x))$ 

Examples of constraints:

- physical restrictions: positiveness, equilibrium laws, ...
- hard goals: target points, production objectives, ...
- safety issues: "staying on the road", collision avoidance, ...



## Broad view on optimisation: Cost function

f allows one to compare the available choices from  $\mathcal{X}$ .

The nature of f (and c) defines several types of problems:

- Linear vs. Nonlinear
- Differentiable vs. Nondifferentiable
- Convex vs. Nonconvex

Examples of objective functions:

- actual cost to minimise: fuel consumption, expenses, ...
- profit to maximise: return on investment, sales, ...

other performance indices: time, distance, velocity, ...
 Typical problem:

$$\min_{c(x) \le 0} f(x) \tag{P_c}$$



## Uncertainty in optimisation: Sources

In some cases, unpredictable events occur during the optimisation process

Examples of sources of uncertainty:

- Power systems, Inventory management: uncertainty on the demand
- Portfolio selection: risks in the share market
- Biological systems: climate, predation
- System stabilization: modelling error, noise

Slide 5/27



## Uncertainty in optimisation: perturbed cost and constraints

Unknown parameter belonging to an uncertainty set:

$$d\in \mathcal{D}$$

It acts on:

- The cost  $\Rightarrow$  F(x, d)
- The constraints  $\Rightarrow C(x, d)$

Main issues:

- Can we make d less uncertain ? Get information (observe) if we can !
- If not, how to write a proper uncertain optimisation problem ? Need to model the uncertainty !



## Observation first vs Decision first

When one is able to observe d then two cases are possible:

► The observation comes *before* the decision (Observation first) ⇒ d is known and we get a parametrized optimization problem:

$$\min_{c_d(x) \le 0} f_d(x) \tag{P_d}$$

with  $c_d(x) = C(x, d)$ ,  $f_d(x) = F(x, d)$  and solutions  $x_d^*$ 

► The observation comes after the decision (Decision first) ⇒ d remains unknown and the problem depends on the representation of the uncertainty

Pitfall to avoid

Consider one value  $d_0$  and work with this one only

Introduction to Stochastic Optimisation

Stochastic Optimisation Study Group

October 2020



## Worst case approach

ldea: consider every situation identically and make a decision based on assuming the worst possible value of  $\boldsymbol{d}$ 

- Cost: f(x) = max<sub>d∈D</sub> F(x, d)
  As if 'nature' is playing against us and we have to assume its best response
- ► Admissible set:  $\mathcal{X} = \{x : \max_{d \in D} C(x, d) \le 0\}$ The constraint must be satisfied for any value of d $\Rightarrow$ robust formulation

Remarks:

- $\blacktriangleright$  The uncertainty is here characterized by  ${\cal D}$
- Very safe but very restrictive so usually not very efficient
  *⇒ conservative* approach

How to avoid conservativeness? Stochastic approach



## Stochastic optimisation

General Idea: consider the various situations differently based on their probability of occurrence

 $\Rightarrow$  Represent the uncertainty as a random variable D valued in  $\mathcal D$ 



## Stochastic optimisation: Expected cost

Idea: make a decision based on the average cost, on a aggregation of the scenarios

- ▶ Cost:  $f(x) = \mathbb{E}[F(x, D)]$
- Admissible set:  $\mathcal{X} = \{x : c(x) \leq 0\}$

- The uncertainty is characterized here by the probability distribution of D
- Typically less conservative but less safe as one does not care about particular realizations of D
- f keeps many properties of F (linearity, convexity, differentiability,...)



## Stochastic optimisation: Constraint in expectation

Idea: put a restriction on the average value of the constraint

 $\blacktriangleright$  Cost: f(x)

Admissible set:  $\mathcal{X} = \{x : c(x) := \mathbb{E}[C(x, D)] \le 0\}$ 

- safer than just a expected cost but not very intuitive
- c keeps many properties of C (linearity, convexity, differentiability,...)



## Stochastic optimisation: Chance constraint

Idea: put a constraint on the probability of admissibility

- $\blacktriangleright$  Cost: f(x)
- Admissible set:  $\mathcal{X} = \{x : \mathbb{P}(C(x, D) \le 0) \ge 1 \alpha\}$ with  $\alpha > 0$  being the level of risk

- Stochastic equivalent of the robust formulation
- $\blacktriangleright$  X becomes a complicated set



## Problems with recourse

What if we could make two decisions, one before observing  $\boldsymbol{D}$  and one after ?

- Choose  $x_0 \in \mathcal{X}_0$  and pay  $f_0(x_0)$
- Observe D
- ▶ Choose  $X_1 = \pi_1(D) \in \mathcal{X}_1$  and pay  $f_1(x_0, X_1, D)$

We get the following two-stage problem:

$$\min_{x_0 \in \mathcal{X}_0, X_1 = \pi_1(D) \in \mathcal{X}_1,} \mathbb{E}[f_0(x_0) + f_1(x_0, X_1, D)]$$

- Typically infinite dimensional problem even if X<sub>0</sub> and X<sub>1</sub> are of finite dimension
- Usually solved by backward decomposition
- Can be extended to multi-stage problem



## Numerical methods

## General remarks

- One can combine the formulations !
- They are versions of (P<sub>c</sub>) with particular structures where f and c are hard to compute

#### Examples of numerical methods

- Expected cost: Stochastic average approximation (SAA), Stochastic gradient methods, etc . . .
- Constraint in expectation: dual methods, primal-dual methods, etc ...
- Chance constraint: deterministic reformulation or approximation, scenario approach, etc ...
- Multi-stage problems: Dynamic Programming, Scenario trees, Reinforcement learning, etc ...



#### Inventory: The newsvendor problem

- A company decides about ordering quantity x of a certain product to satisfy demand d (unknown at purchase time)
- Cost per unit is c
- ▶ If d > x, additional order at unit price  $b \ge 0$  is made
- $\blacktriangleright$  b > c, i.e. backorder penalty cost is larger than ordering
- ▶ If d < x, a holding cost of h(d x) is incurred

#### Total cost

$$F(x,d) = cx + b \max\{d - x, 0\} + h \max\{x - d, 0\}$$

#### Objective

Minimise the total cost F(x, d), where x is the decision variable and d is a parameter.

Introduction to Stochastic Optimisation | Stochastic Optimisation Study Group | October 2020 |

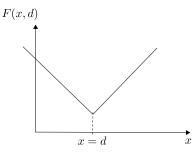


## Cost function

We can re-write the objective function as

$$F(x,d) = \max\{(c-b)x + bd, (c+h)x - hd\},\$$

which is a piecewise linear function with minimum attained at  $\bar{x} = d$ , provided the demand is **known!** 



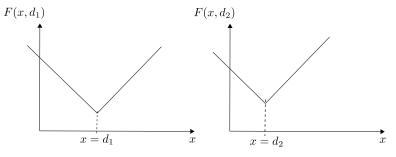


## What should we do when d is unknown when ordering?

The ordering decision should be made **before** a realisation of the demand becomes known. View demand D as a random variable (d is a particular realisation of D)

Silly attempt: Optimise  $\min_{x\geq 0} F(x, D)$ 

> This problem does not make sense! F(x, D) is random



One x cannot simultaneously optimise both functions!



## Attempt 2: Plan for the average case

- ▶ Suppose the cdf of *D* is given by  $H(x) := \mathbb{P}(D \le x)$
- Let  $\mu \coloneqq \mathbb{E}[D]$
- ▶ Then,  $\arg \min_{x \ge 0} F(x, \mu) = \mu$
- The optimal policy is to order a quantity  $\mu$

#### Remark

This approach might be more sensitive to random perturbations of the empirical data.



#### Attempt 3: Plan for worst case

- Suppose that we know upper and lower bounds on the demand d, i.e. ℓ ≤ d ≤ u
- Consider  $\min_{x \ge 0} \{ f(x) \coloneqq \max_{d \in [\ell, u]} F(x, d) \}$
- While making a decision x, one is prepared for the worst possible outcome of the maximal cost
- The optimal solution is

$$x^* = \frac{h\ell + bu}{h+b}$$

- ▶ If the holding cost is zero (h = 0), then the optimal solution is obviously  $x^* = u$  (always order the maximum demand)
- Often a conservative approach!



## Proof

▶ Recall that  $F(x, d) = \max\{(c - b)x + bd, (c + h)x - hd\}$ , then  $\max_{d \in [\ell, u]} F(x, d) = \max\{F(x, \ell), F(x, u)\}$ 

 $\blacktriangleright$  We should look at the problem in  $x \in [\ell, u]$  , then

$$\min_{x \ge 0} \max_{d \in [\ell, u]} F(x, d)$$
  
=  $\min_{x \in [\ell, u]} \max\{F(x, \ell), F(x, u)\}$   
=  $\min_{x \in [\ell, u]} \max\{cx + h[x - \ell]_+, cx + b[u - x]_+\}$ 

► This is a piecewise linear convex function, and the optimal solution is attained at the point where h(x - ℓ) = b(u - x), completing the proof.



## A suitable framework

- $\blacktriangleright$  The probability distribution of D is known
- ▶ It makes sense to talk about the expected value  $\mathbb{E}[F(x,D)]$  of the total cost as a function of x
- Consequently,

$$\min_{x \ge 0} \{ f(x) \coloneqq \mathbb{E} \left[ F(x, D) \right] \}$$

Slide 21/27



## Solution to $\min_{x\geq 0} \mathbb{E}\left[F(x,D)\right]$

- This particular example can be solved in closed form
- $\blacktriangleright \text{ Recall } H(x) \coloneqq \mathbb{P}\{D \leq x\}, \ H(x) = 0 \text{ for all } x < 0$
- The optimal solution is equal to the quantile

$$x^* = H^{-1}\left(\frac{b-c}{b+h}\right)$$

It is not always this easy...

- The newsvendor problem is about the only stochastic program that admits a simple "closed-form" solution
  - In general, we must solve instances **numerically**!



## Proof

#### Recall that

$$F(x, D) = cx + b \max\{D - x, 0\} + h \max\{x - D, 0\}$$
  
= max{(c - b)x + bD, (c + h)x - hD}

Moreover,

$$f'(z) = \frac{\partial}{\partial z} \mathbb{E} \left[ cz + b \max\{D - z, 0\} + h \max\{z - D, 0\} \right]$$
$$= c - b\mathbb{P}[D \ge z] + h\mathbb{P}[D \le z]$$
$$= c - b(1 - H(z)) + hH(z)$$
$$= c - b + (b + h)H(z)$$



## Proof

Then,

$$\mathbb{E}[F(x,D)] = b\mathbb{E}[D] + (c-b)x + (b+h)\int_0^x H(z)dz$$

To obtain the minimum we take the derivative of the right-hand side and equate it to zero. Consequently, the optimal solutions satisfy

$$(b+h)H(x^*) + c - b = 0$$

and the result follows.



## Numerical comparison

$$F(x,D) = cx + b[D - x]_{+} + h[x - D]_{+}$$

• Cost per unit: 
$$c = 1$$

- **Re-ordering cost:** b = 5
- **•** Holding cost: h = 0.5
- **Demand:** Normally distributed  $\mathcal{N}(100, 14^2)$

Plan for	Optimal solution $x^*$	$\mathbb{E}[F(x^*, D)]$
Average demand	100	\$ 131
Worst case (58 $\leq D \leq 142$ )	134	\$ 152
Expected cost	112	\$ 126



## Chance constraints

#### Remark

For particular realisations of the demand,  $F(x^*, D)$  can be quite different from  $\mathbb{E}[F(x^*, D)]$ . A natural question is whether we can control the risk of the cost F(x, D) to be not "too high".

- For a given threshold  $\tau > 0$ , we want  $F(x, D) \le \tau$  to be satisfied for all possible realisations of  $D \in \mathcal{D}$
- Can be quite restrictive if the uncertainty set D is large!
- We thus introduce a constraint

$$\mathbb{P}\{F(x,D) \le \tau\} \ge 1 - \alpha$$

For  $x \leq \tau/c$ , this becomes

$$H\left(\frac{(b-c)x+\tau}{b}\right) - H\left(\frac{(c+h)x-\tau}{h}\right) \ge 1 - \alpha$$



## Newsvendor problem

The newsvendor is a simple example of a two-stage problem

- First stage: Before a realisation of D is known, one makes a decision about x
- Second stage:
  - After a realisation d of D becomes known, it may happen that d > x (Random "stuff" happens)
  - ► The company then takes a recourse action of ordering d x at a higher cost of b > c (We attempt to repair the havoc)

Slide 27/27

The evolution of information is of paramount importance!

## THANK YOU FOR YOUR ATTENTION!