

Errata Listing for  
*Compression and Coding Algorithms*,  
Kluwer Academic Publishers, 2002

Alistair Moffat and Andrew Turpin

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**Preface**

**Chapter 1**

*Page 9*

The claim on disk transfer rates is somewhat dated, and modern disks operate at rates in excess of 50 MB/second. Seek times are still in the 8–10 millisecond range.

[Shingo Mamada, September 2002]

**Chapter 2**

**Chapter 3**

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For example, when  $b = 5$  the codeword lengths

*should be replaced by*

For example, when  $n = 5$  the codeword lengths

[Ke Shi, August 2016]

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For example, a Rice code with  $k = 1$  assigns the codewords “00”, “01”, “100”, “101”, “1100”, and “1101”, to symbols 0, +1, -1, +2, and -2, respectively, and is biased in favor of the positive values.

*should be replaced by*

For example, a Rice code with  $k = 1$  assigns the codewords “00”, “01”, “100”, “101”, “1100”, and “1101”, to symbols 0, +1, -1, +2, -2, and +3, respectively, and is biased in favor of the positive values.

[Eduardo Morras, December 2001]

This error reflects exactly the point being made in that paragraph – that an asymmetric set of codeword lengths is inappropriate.

**Page 43**

In line 1 of *centered\_minimal\_binary\_encode* it should say “set  $long \leftarrow 2 \times n - 2^{\lceil \log_2 n \rceil}$ ”, that is, no reference to  $x$  while the number of *long* codewords is being calculated.

[André Osterhues, May 2004]

**Chapter 4**

**Page 64**

At the bottom of the page we mention a decoding mechanism by Hashemian [1995], that uses  $k$ -bit tables to provide either the decoding and number of bits for a code if that length is less than or equal to  $k$ , or a pointer to another table for codes longer than  $k$  bits. This decoding method was previously used in the inflate routines in Info-ZIP’s unzip and gzip, both developed in 1992 by Mark Adler.

[Mark Adler, October 2004]

**Page 70–77**

The mechanism described in this section is the work of Moffat and Turpin (1998). Somehow or other we missed citing our own paper!

[Alistair Moffat, December 2001]

**Page 86**

The first assignment (to *sym*) in step 10 of Algorithm 4.7 is not required. Then, in the second assignment in step 11, the right-hand-side expression should be  $w[\text{code\_len}] - 1$ , not the “+1” shown in the book.

[Andrew Turpin, October 2003]

**Chapter 5**

**Page 108**

In Table 5.3, when coding the sixth symbol (which is a 4), the value of  $r$  that is given is incorrect, it should be 6 rather than the 4 that is shown (calculated as  $60/10$ ).

[Wendy Cameron, April 2004]

**Chapter 6**

**Chapter 7**

**Chapter 8**

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Another theme that has been explored by a number researchers  
*should be replaced by*

Another theme that has been explored by a number of researchers

**Page 241**

The inversion coding technique of Arnavut (2000) should also have been mentioned as an alternative to the MTF transformation. It is an alternative way of dealing with the localized nature of the post-BWT string. Interpolative coding (described in Chapter 3) can also be used in a similar way, as described by Moffat and Stuiver (2000).

[Alistair Moffat, December 2001]

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Another recent paper examining searching over compressed text is by Navarro et al. (2001).

[Alistair Moffat, December 2001]

## **Chapter 9**

**Page 232**

A recent PPM implementation of Dmitry Shkarin (2002) obtains extremely good compression, with one version reported to attain a “Calgary average” of 1.92 bits per character. A link to an executable is provided at <http://datacompression.info/PPM.shtml>.

[Sebastian Deorowicz, August 2002]

**Page 241**

The blocksize for SZIP is 4.1 MB, not 4.3 MB as listed.

[Sebastian Deorowicz, August 2002]

**Page 241**

Another recency transformation that can be used with the BWT is “Distance Coding”. Deorowicz (2002) includes a description of distance coding in his recent paper. Elias (1987) also considers various recency transformations.

[Sebastian Deorowicz, August 2002]

## **Bibliography**

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The third paper listed is by “Guazzo”, not “Guauzzo”.

[Sebastian Deorowicz, August 2002]

**Page 263**

The paper Liddell and Moffat (2002) will appear in the proceedings of the 2002 *IEEE Data Compression Conference*.

[Alistair Moffat, December 2001]

**Page 265**

The seventh paper listed is by “Pasco”, not “Pascoe”.

[Sebastian Deorowicz, August 2002]

## Index

## References

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