

# **The $V^*$ -Diagram: A Query-Dependent Approach to Moving KNN Queries**

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# Motivation

Consider two scenarios:

- a driver in a GPS-equipped car finding the nearest gas station along the route of a trip;
- a tourist walking in the city looking for the nearest ATM.

These scenarios are examples of *moving  $k$  nearest neighbor queries (MkNN)*.

# Simple Approach

## The Voronoi Diagram

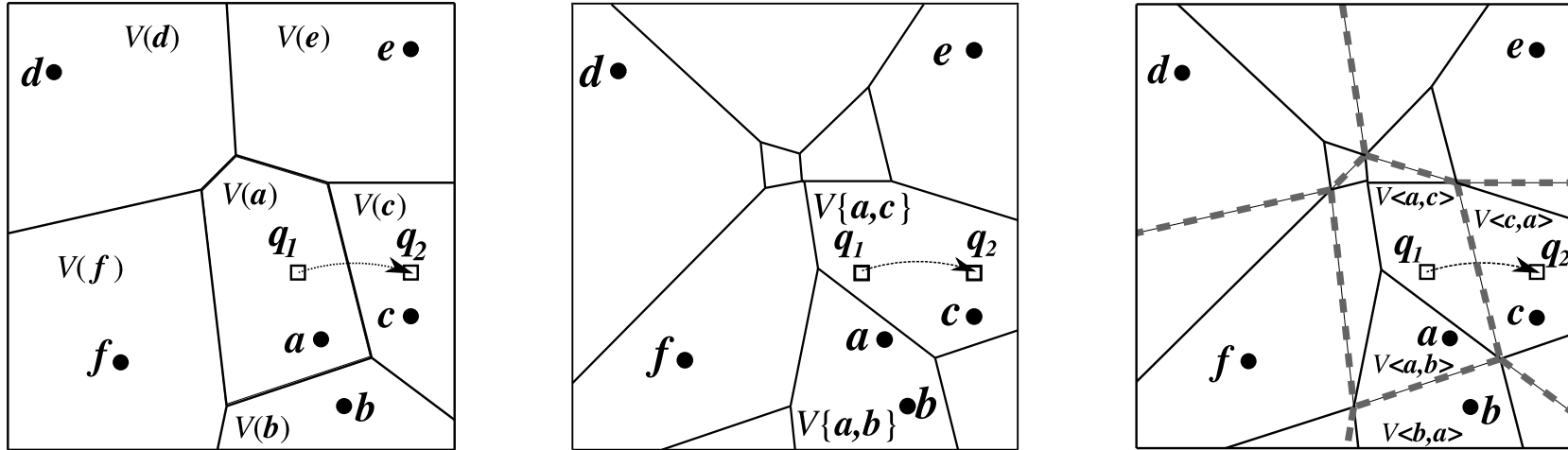


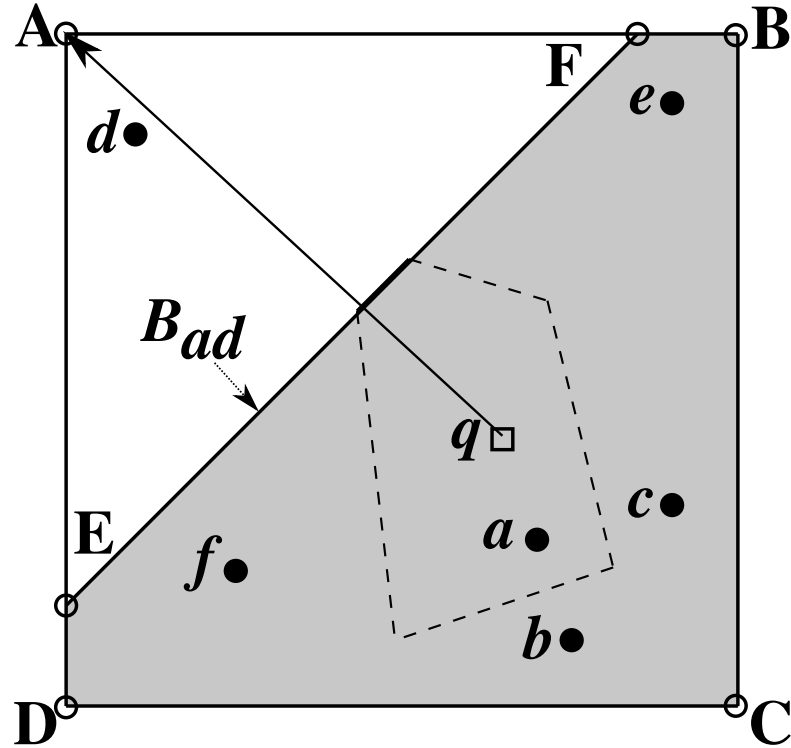
Figure 1: Voronoi diagrams

### Drawbacks:

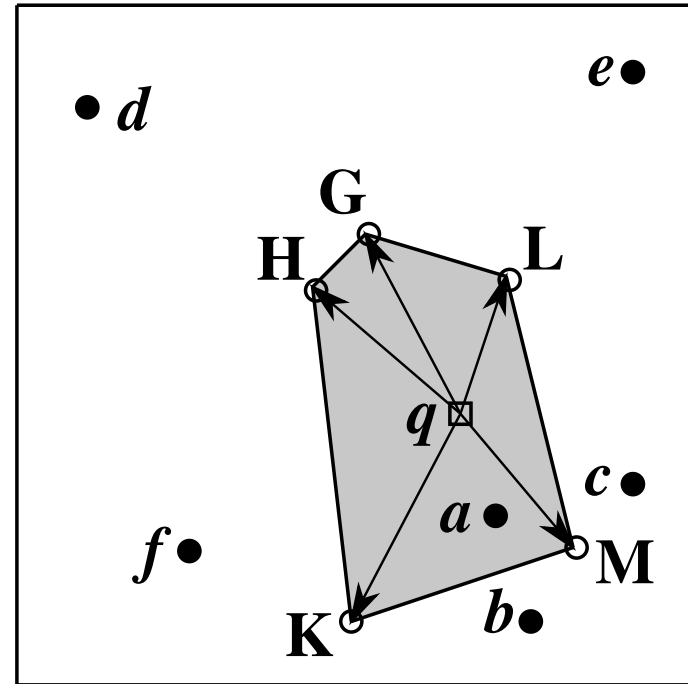
1. Expensive precomputations
2. Inefficient update operations
3. No support for dynamically changing  $k$  values

# Best Existing Approach

Influence-set Retrieval [Zhang et al., 2003]



(a) Bisector  $B_{ad}$  is discovered as a boundary.



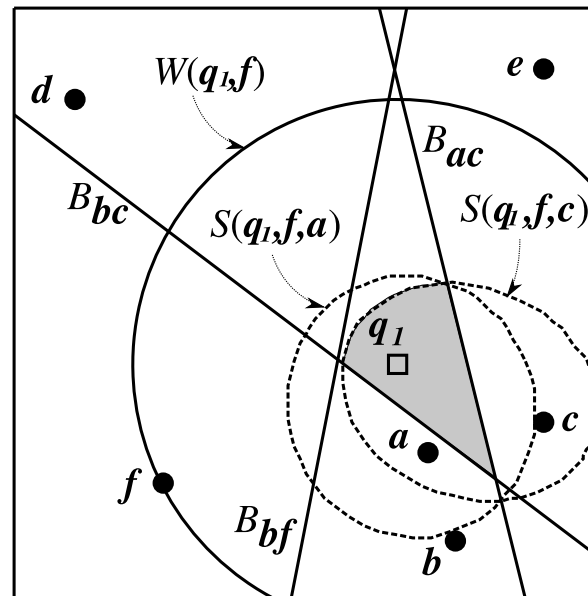
(b) All boundaries are discovered

Figure 2: Computing a Voronoi cell locally

# Our Approach: $V^*$ -Diagram

Objectives:

1. Requires *no precomputation*
2. Supports dynamic *insertions/deletions* of objects
3. Handles *dynamically changing  $k$*

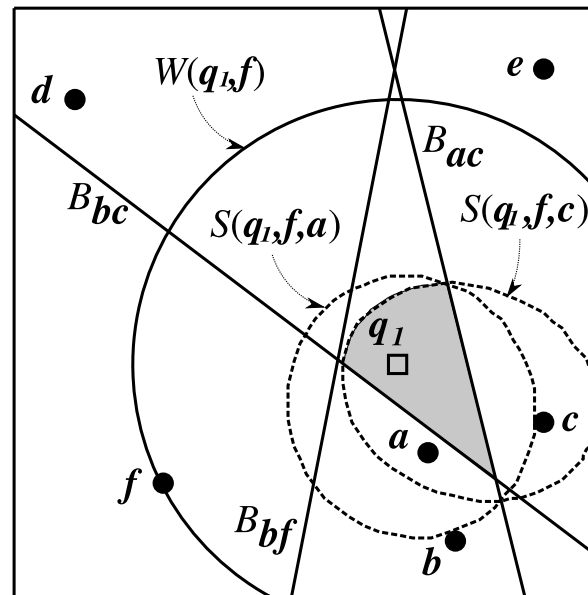


# Our Approach: $V^*$ -Diagram

Objectives:

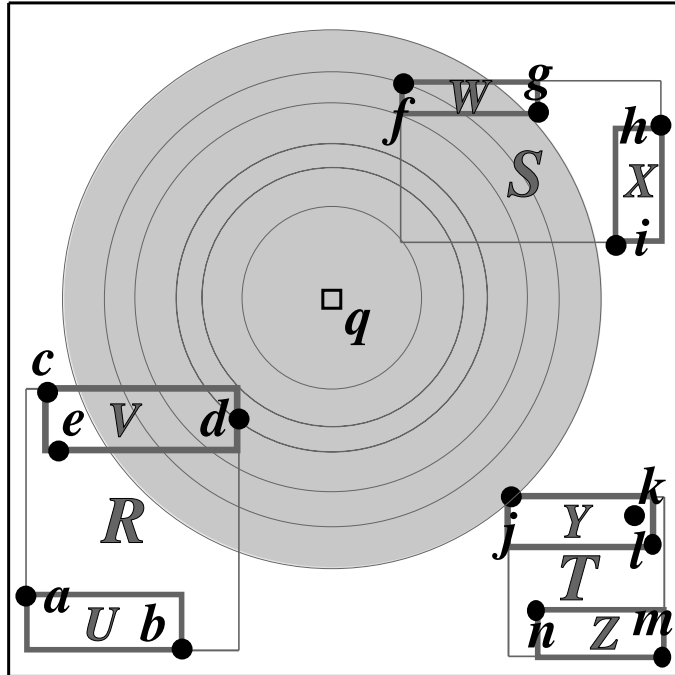
1. Requires *no precomputation*
2. Supports dynamic *insertions/deletions* of objects
3. Handles *dynamically changing  $k$*

Result: Outperforms the best practice [Zhang et al.]  
by *2 orders of magnitude*



# The $V^*$ -Diagram

## Known Region



If the known NNs to  $q$  are

$$\{d, f, j\},$$

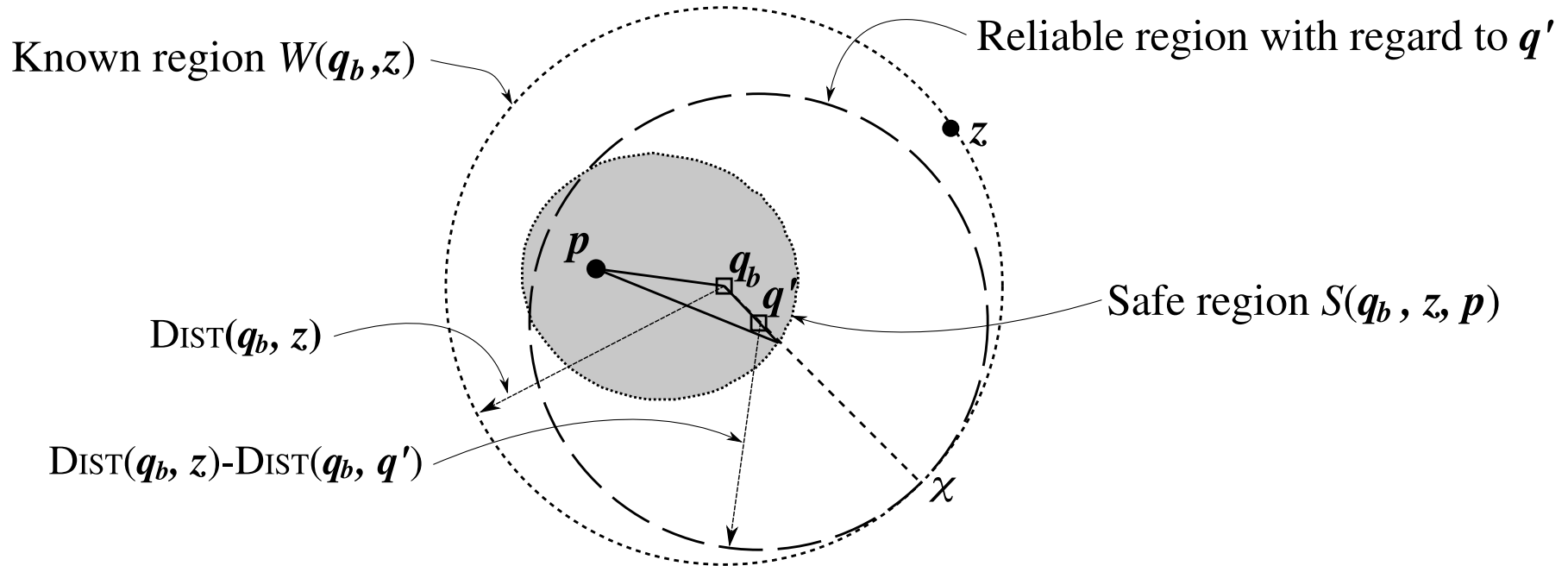
the know region  $W(q, j)$  is

$$\{v : dist(q, v) \leq dist(q, j)\}.$$

# The $V^*$ -Diagram

## Safe region wrt a data point

We retrieve  $(k + x)$  objects. In this example,  $k$  and  $x$  are 1, so we retrieve  $p$  and  $z$ .



If  $q' \in S(q_b, z, p)$  then,

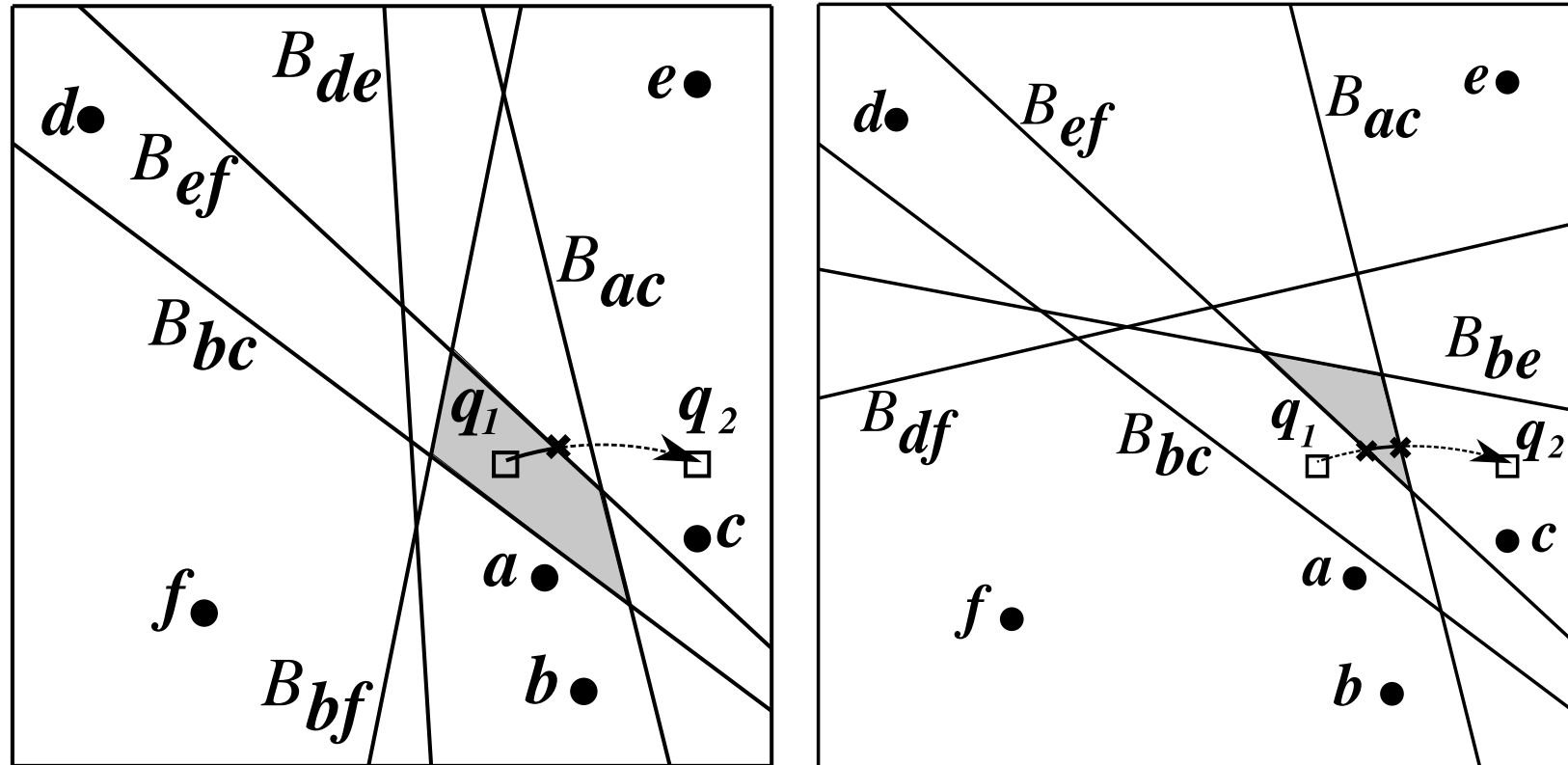
$$\forall p' \notin W(q_b, z), dist(q', p) < dist(q', p').$$

$$S(q_b, z, p) = \{q' : dist(p, q') \leq dist(q_b, z) - dist(q_b, q')\}$$



# The $V^*$ -Diagram

The Fixed-rank Region (FRR) [Kulik and Tanin, 2006]



(a)  $\langle a, c, b, f, e, d \rangle$

(b)  $\langle a, c, b, e, f, d \rangle$

Figure 3: Incremental rank update

# The $V^*$ -Diagram

## Integrated Safe Region (ISR) and $V^*$ - $k$ NN

ISR is an intersection of

1. the safe region wrt  $k^{th}$  NN,  $S(\mathbf{q}_b, \mathbf{z}, \mathbf{p}_k)$ ;
2. the FRR of the  $(k+x)$  NNs of  $\mathbf{q}_b$ .

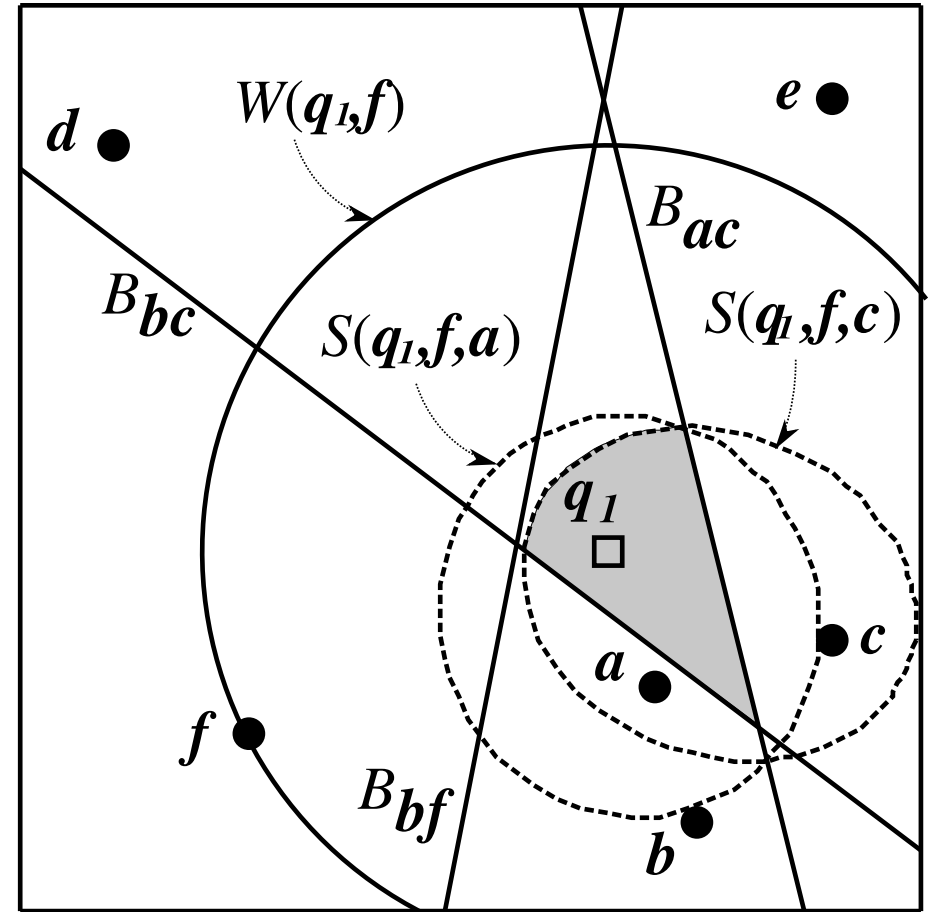


Figure 4:  $V^*$ - $k$ NN Example ( $k = 2, x = 2$ )

# V\*-kNN Algorithm

<http://www.csse.unimelb.edu.au/~sarana/demo.html>

Legend & Events  
Java Applet Window

- Known Region: The area that has been searched
- Reliable Region: The region that contains reliable\* object w.r.t.  $q$
- Bisector of two data points
- Safe Region (SR) wrt  $k$ -th NN: Region that keeps the  $k$  NNs reliable\* w.r.t.  $q$
- Fixed-Rank Region (FRR): The region that the ranking
- Integrated Safe Region (ISR): A  $k$ NN safe region
- Known Object: One of the  $(k+x)$  NNs
- Unknown Object: An object yet to be discovered

$i:S_j$   $S_j$  is the  $i$ -th NN of  $q$

\* An object  $p$  is reliable w.r.t.  $q$  if it is guaranteed that there is no object outside the known region is nearer to  $q$  than  $p$

Instruction  
Animation complete. Now the query point can be moved around by dragging it

Historical Events

Current Event

Applet Viewer: Demo.class

Applet

### Demonstration of the V\*-Diagram and V\*-kNN Algorithm

$k = 2$   $x = 3$   SR  FRR  $n = 10$  Start Over Regenerate Data

Authors: Sarana Nutanong, Rui Zhang, Egemen Tanin and Lars Kulik  
Email: sarana@csse.unimelb.edu.au  
Reference: The V\*-Diagram: a query-dependent approach to moving KNN queries. VLDB'08

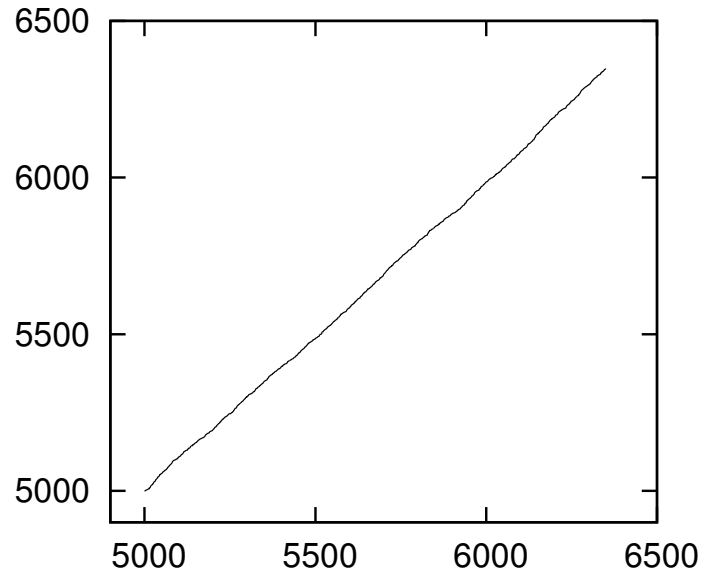
Applet started.

# Experiments

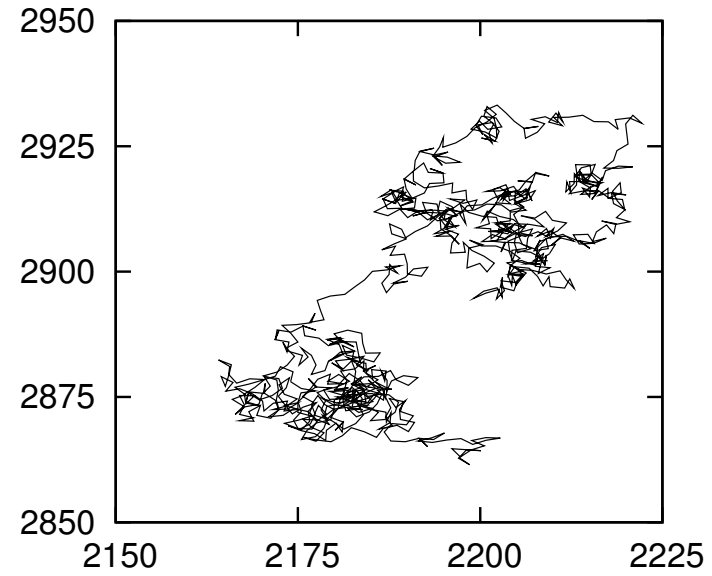
- Data Structure: R\*-trees (1-kB block size).
- Comparative Method: RIS- $k$ NN [Zhang et al.]
- Datasets:
  - (U) 25,000 of data points in uniform distribution
  - (Z) 25,000 of data points in Zipfian distribution
  - (C) 65,743 postal addresses from California
  - (N) 119,897 postal addresses from North-Eastern USA

# Experiments

## Trajectories



(a) Directional (D)

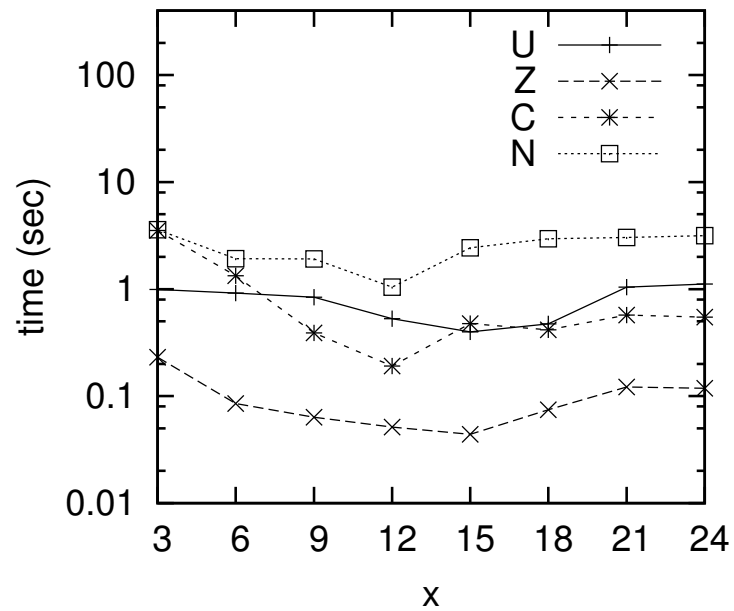


(b) Random (R)

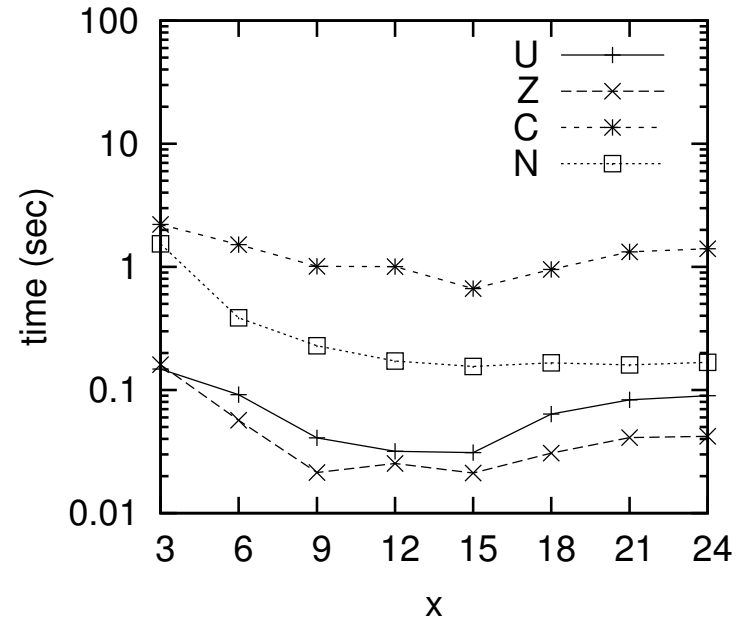
Figure 5: Trajectory types

# Experiments

total cost wrt  $x$



(a) Total cost (D)

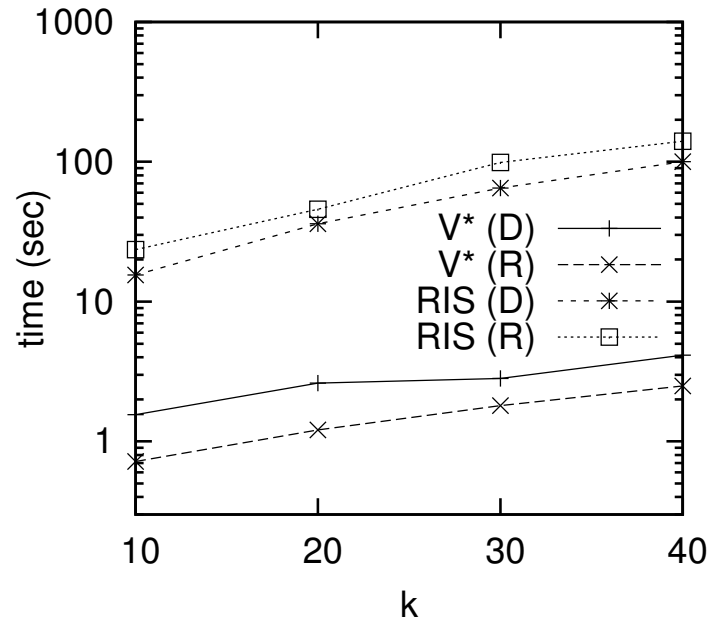


(b) Page access (D)

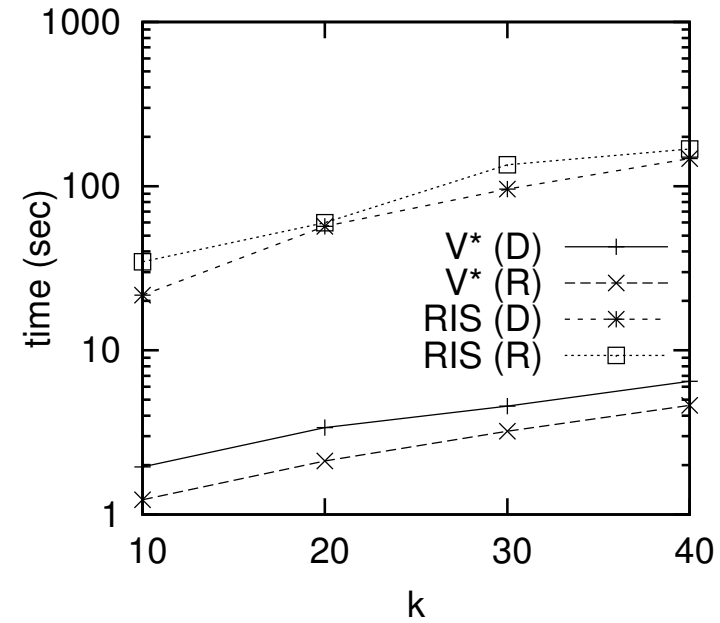
Figure 6: Effect of  $x$

# Experiments

total cost wrt  $k$



(a) Total Cost (California)

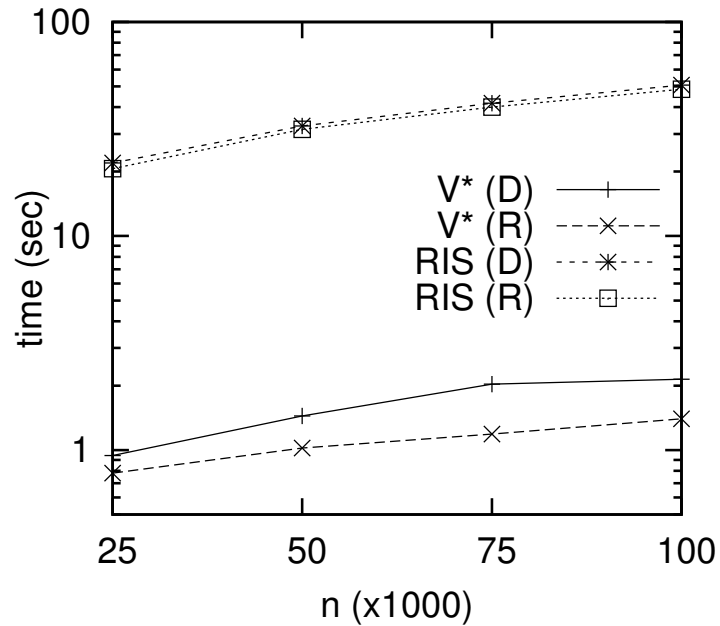


(b) Total Cost (North-Eastern USA)

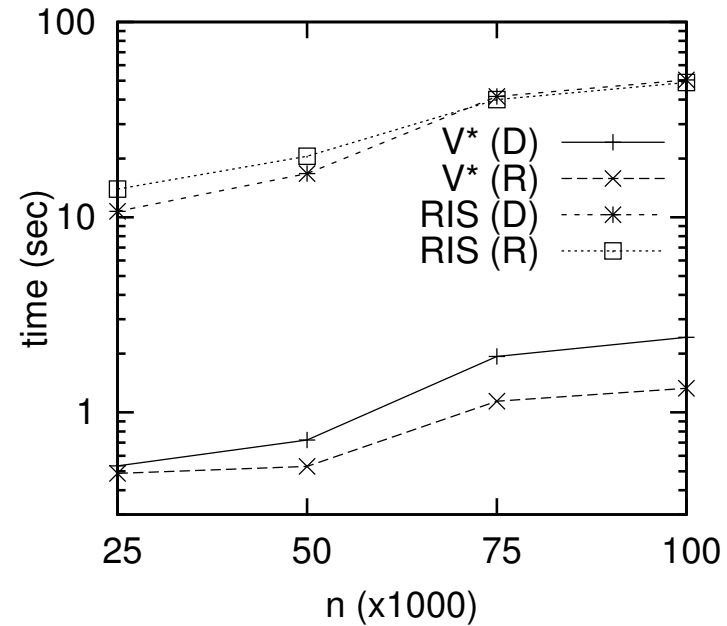
Figure 7: Effect of  $k$

# Experiments

total cost wrt  $n$



(a) Total Cost (Uniform)



(b) Total Cost (Zipfian)

Figure 8: Effect of dataset size



# Cost model

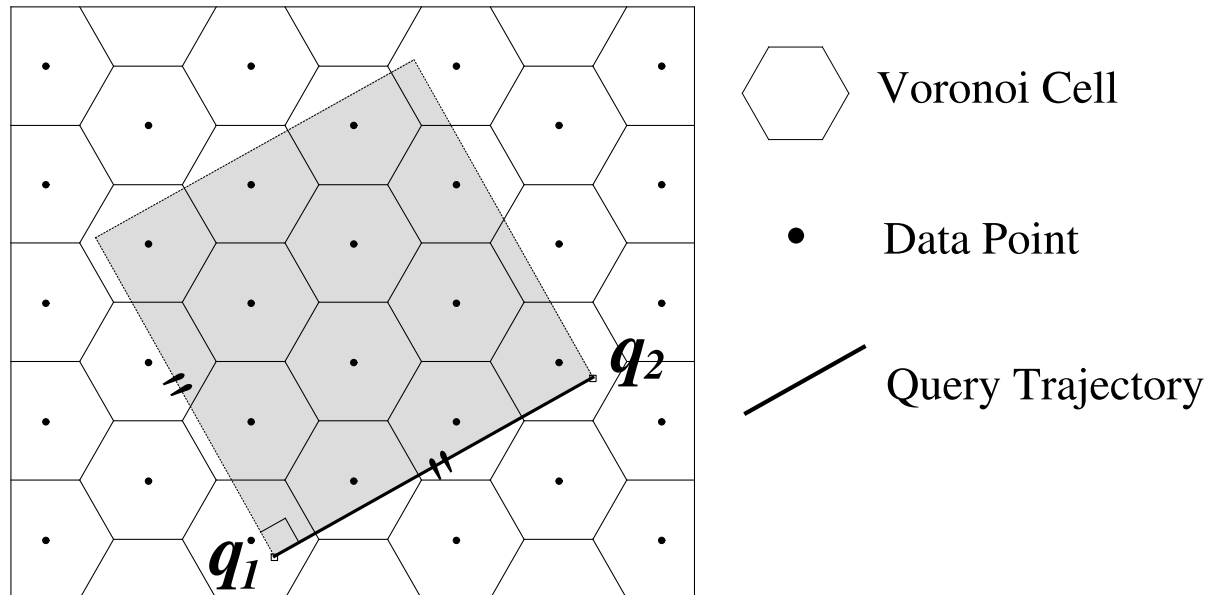
## RIS- $k$ NN

The number of the  $k$ VD cells in 2D space is approximated as

$$2kn \text{ [Okabe et al., 1992].}$$

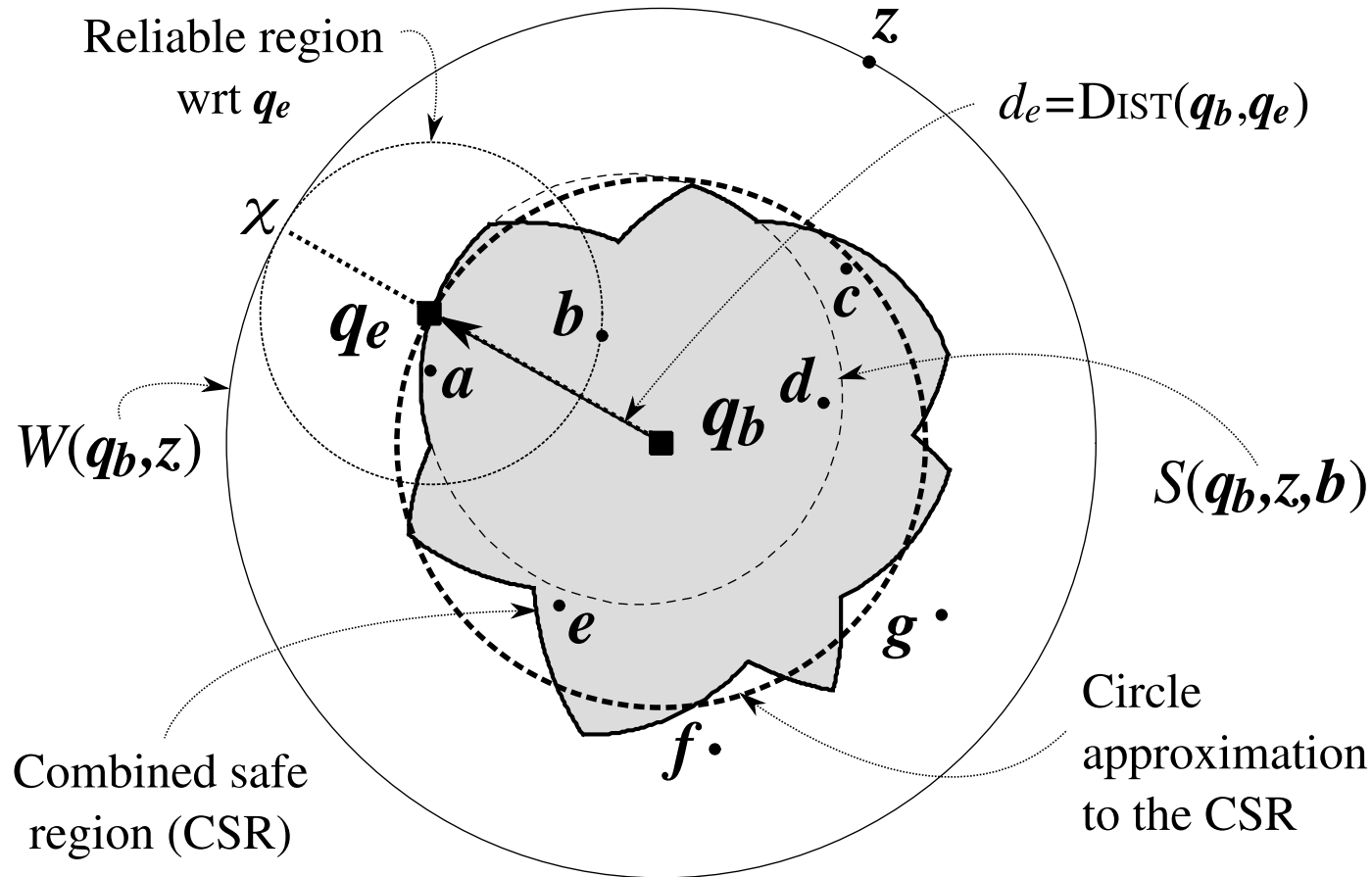
For a given trajectory length  $l$ , the number  $n_v$  of  $k$ VD cells crossed by the trajectory is given by

$$n_v = l\sqrt{2kn}.$$



# Cost model

$V^*$ - $k$ NN

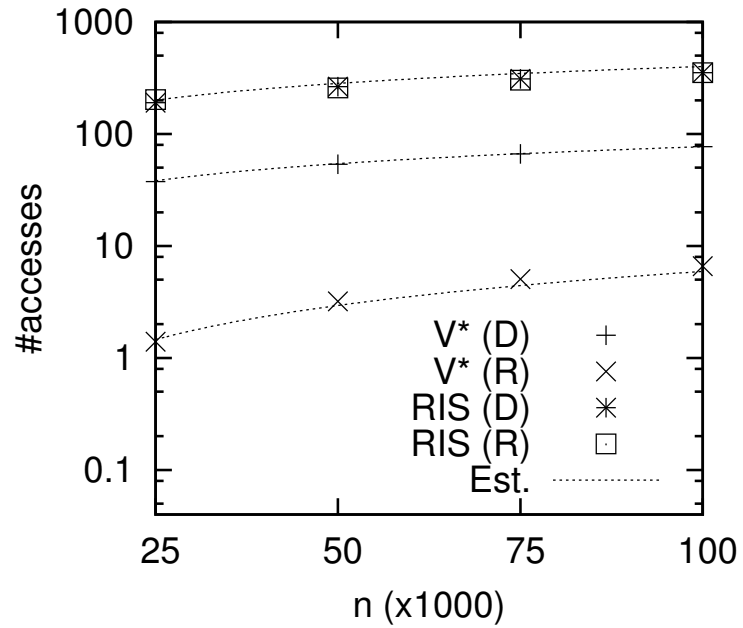


Directional:  $n_b = l/d_e$ .

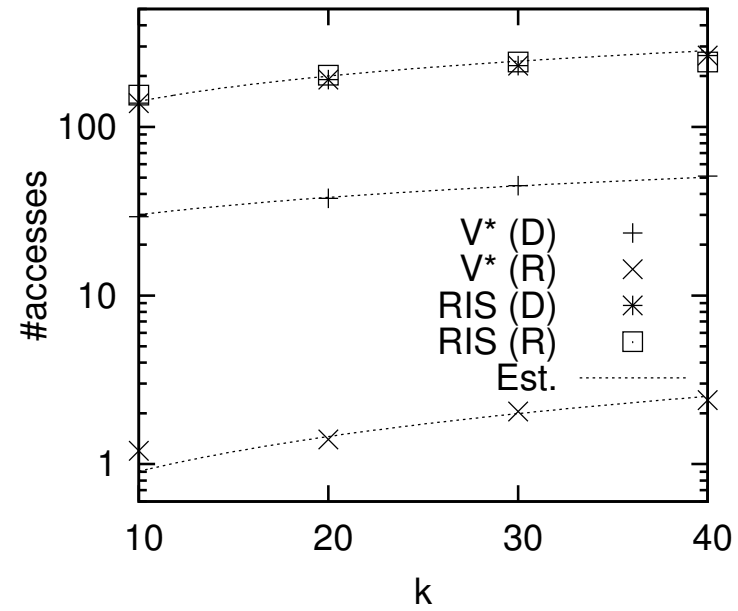
Random:  $n_b = ls/d_e^2$ , where  $s$  is the step size.

# Experiments

## Cost Model



(a) Effect of  $n$



(b) Effect of  $k$

Figure 9: Cost model validation

# The $V^*$ -Diagram in a spatial network

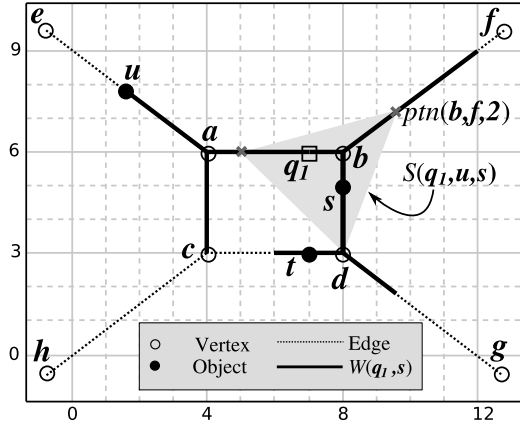


Figure 10: Safe region

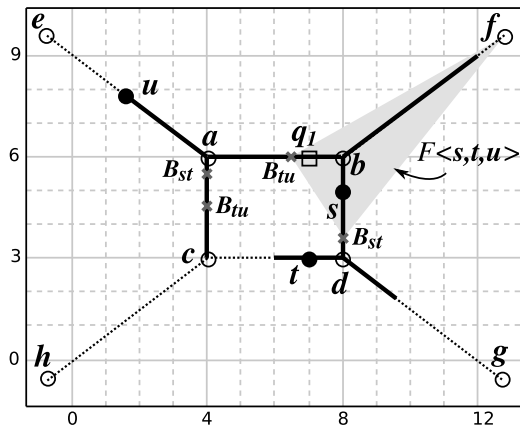


Figure 11: Fixed-rank region

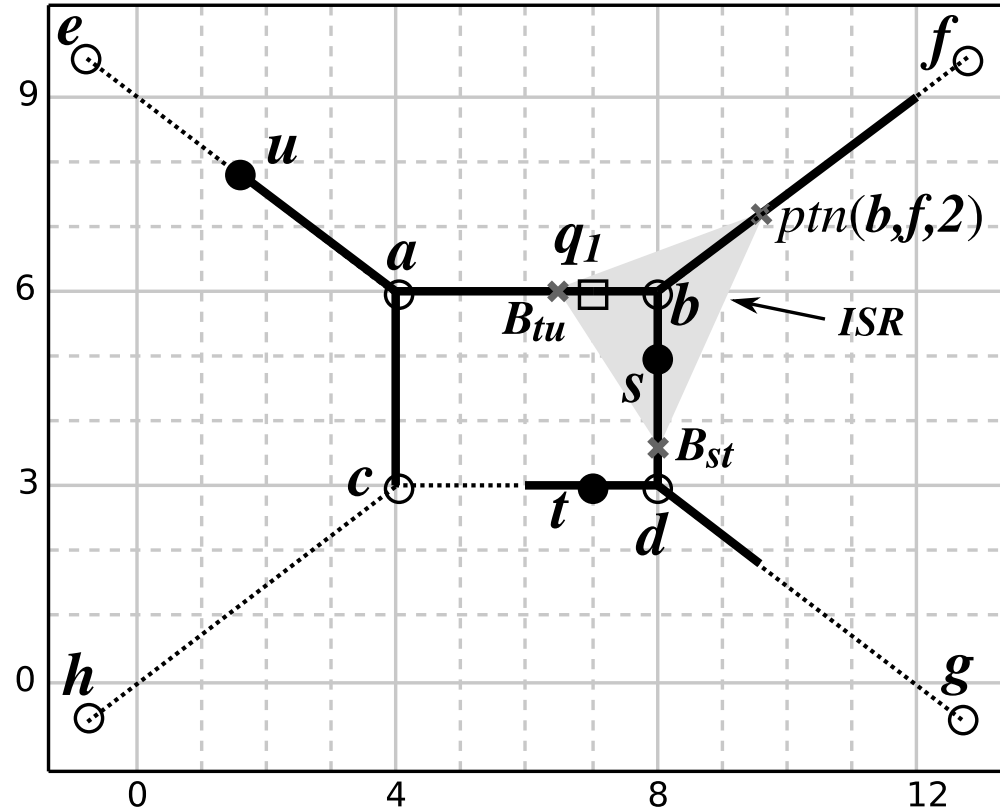


Figure 12: ISR is  $S(q_1, u, s) \cap F\langle s, t, u \rangle$

# Experiments

## The $V^*$ -Diagram in a spatial network

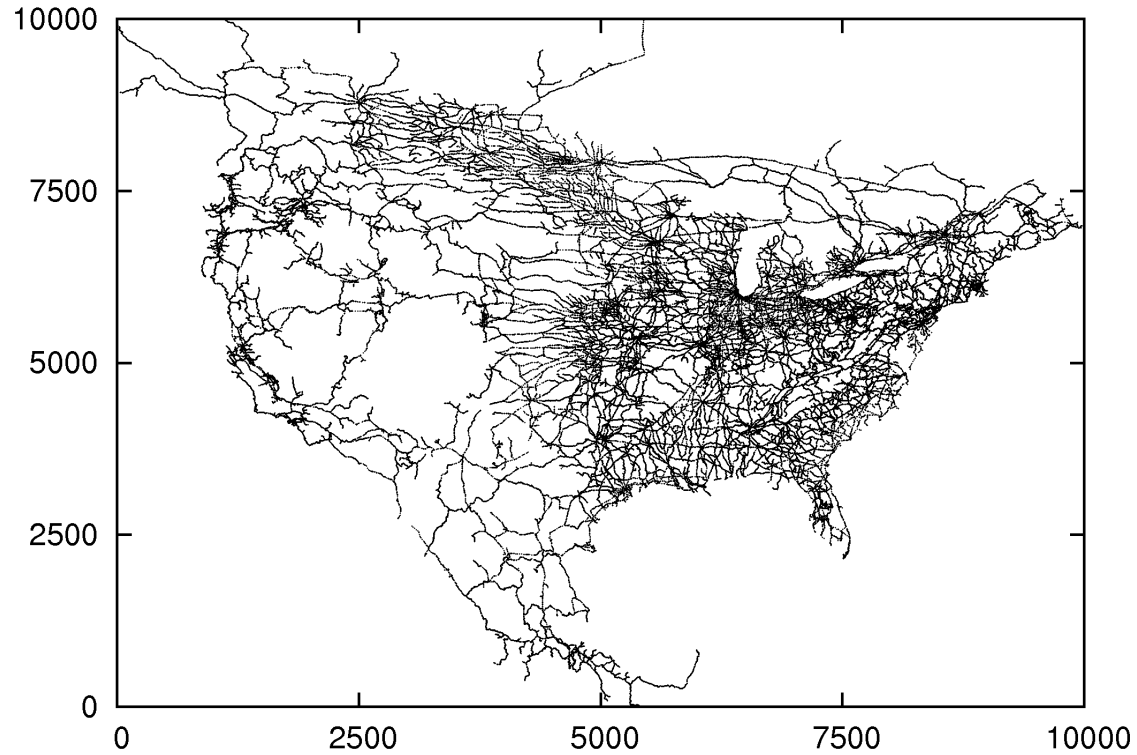
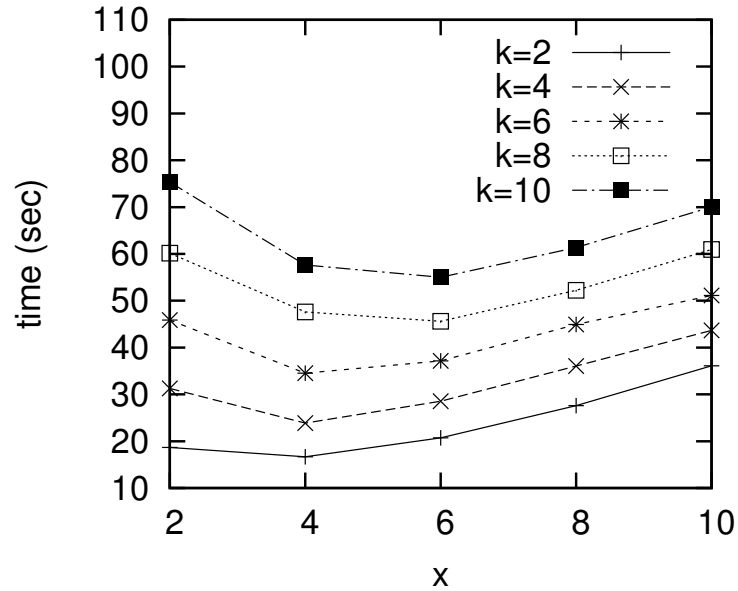


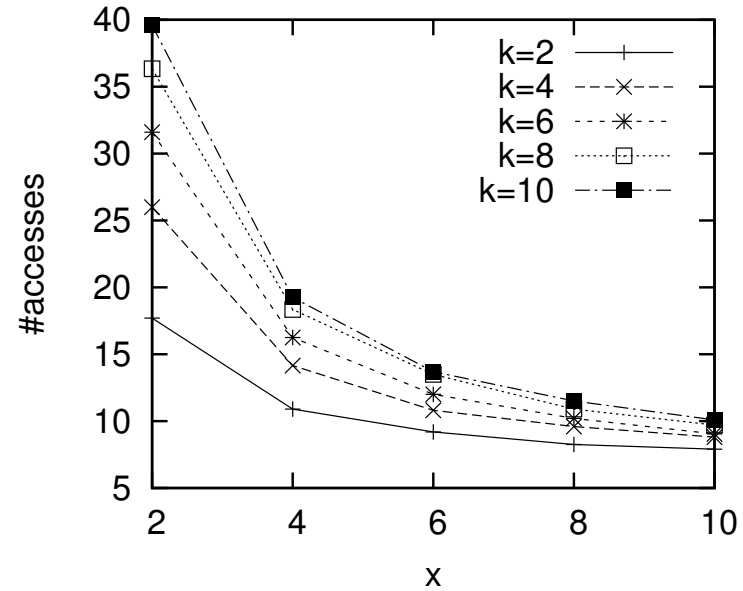
Figure 13: Road network in north America (175,813 nodes and 179,179 edges)

# Experiments

## The $V^*$ -Diagram in a spatial network



(a) Total Response Time

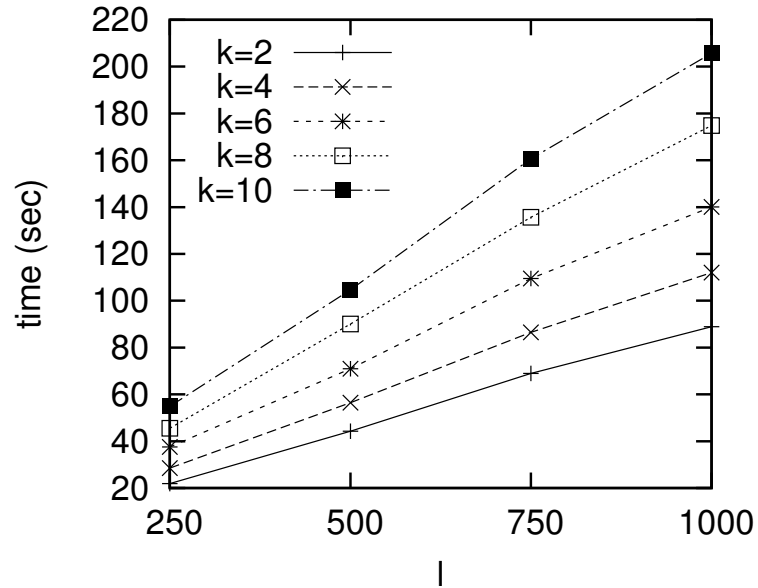


(b) Access Cost

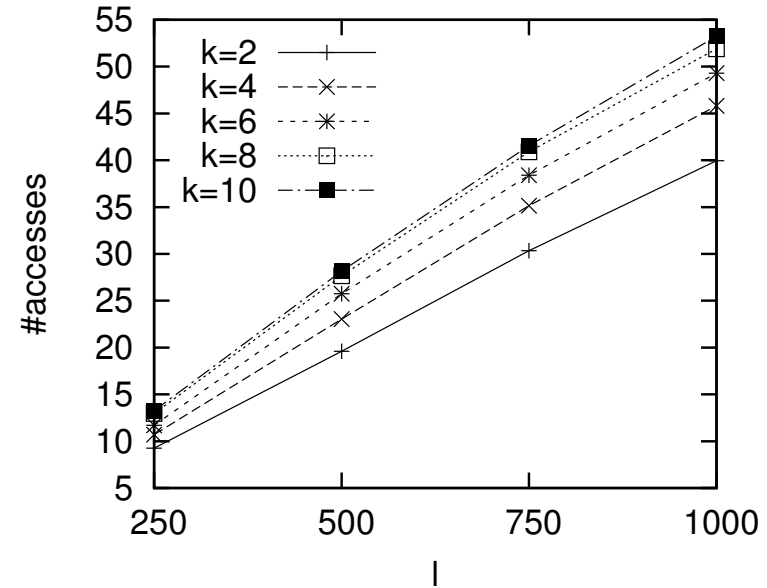
Figure 14: Spatial network: effect of  $x$

# Experiments

## The $V^*$ -Diagram in a spatial network



(a) Total Response Time



(b) Access Cost

Figure 15: Spatial network: effect of  $l$

# Conclusions

- The  $V^*$ -Diagram constructs a safe region using:
  1. the location of the query point,
  2.  $k$ NN-search coverage (known region),
  3. known data points.
- $V^*$ - $k$ NN is *local*, *incremental* and *dynamic*.
- $V^*$ - $k$ NN outperforms the best existing technique by two orders of magnitude.
- The  $V^*$ -diagram is a general philosophy, which can be applied to most safe region based techniques.



# Related Publications

- S. Nutanong, R. Zhang, E. Tanin, L. Kulik: Analysis and Evaluation of  $V^*$ - $k$ NN: An Efficient Algorithm for Moving  $k$  Nearest Neighbor Queries. To appear in VLDB Journal.
- S. Nutanong, R. Zhang, E. Tanin, L. Kulik:  $V^*$ - $k$ NN: An Efficient Algorithm for Moving  $k$  Nearest Neighbor Queries (Demo). ICDE 2009: 1519-1522.
- S. Nutanong, R. Zhang, E. Tanin, L. Kulik: The  $V^*$ -Diagram: a query-dependent approach to moving KNN queries. PVLDB 1(1): 1095-1106 (2008).

# Key References

- Lars Kulik, Egemen Tanin: Incremental Rank Updates for Moving Query Points. GIScience 2006:251-268.
- Atsuyuki Okabe, Berry Boots, Kokichi Sugihara, Sung Nok Chiu: Spatial Tessellations: Concepts and Applications of Voronoi Diagrams. John Wiley & Sons, Inc., 1992.
- Jun Zhang, Manli Zhu, Dimitris Papadias, Yufei Tao, Dik Lun Lee: Location-based Spatial Queries. SIGMOD 2003:443-454.