Analytical results for the multi-objective design of model-predictive control

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Abstract

In model-predictive control (MPC), achieving the best closed-loop performance under a given computational capacity is the underlying design consideration. This paper analyzes the MPC tuning problem with control performance and required computational capacity as competing design objectives. The proposed multi-objective design of MPC (MOD-MPC) approach extends current methods that treat control performance and the computational capacity separately—often with the latter as a fixed constraint—which requires the implementation hardware to be known a priori. The proposed approach focuses on the tuning of structural MPC parameters, namely sampling time and prediction horizon length, to produce a set of optimal choices available to the practitioner. The posed design problem is then analyzed to reveal key properties, including smoothness of the design objectives and parameter bounds, and establish certain validated guarantees. Founded on these properties, necessary and sufficient conditions for an effective and efficient optimizer are presented, leading to a specialized multi-objective optimizer for the MOD-MPC being proposed. Finally, two real-world control problems are used to illustrate the results of the tuning approach and importance of the developed conditions for an effective optimizer of the MOD-MPC problem.

1. Introduction

Model-predictive control (MPC) is a typically computationally expensive method of approaching the control of constrained systems. As a result, the computational capacity required at each sampling instant is a consideration in the overall design process. This is particularly true in systems with fast dynamics, where there is often significant conflict between the complexity of the problem considered at each time step and the available time to find a solution. The close interrelation between control performance and required computational capacity warrants that these indices are analyzed in synchrony to streamline the design process and avoid unnecessary costs. Both objectives depend on a number of tuning parameters of the optimal control problem including, but not limited to, the sampling time, prediction horizon length, and fidelity/order of the prediction model.

Previously, much focus has been given to find the best control performance in a single-objective optimization design problem, separate to the consideration of the required computational capacity. However, there are still a number of knowledge gaps in existing MPC tuning approaches. MPC tuning for control performance is mostly done via methods that rely on rules-of-thumb and general guidelines (Garriga & Soroursh, 2010; Qin & Badgwell, 2003; Rani & Unbehauen, 1997). Further developments have been made consequently, based on metaheuristics such as particle swarm optimization (Júnior, Martins, & Kalid, 2014) and genetic algorithms (van der Lee, Svrcek, & Young, 2008), as well as gradient descent in run-time with imperfect plant model for the single-objective performance optimization of MPC (Bunin, Fraire, François, & Bonvin, 2012).

Several multi-objective optimization approaches for control system design have also been studied for the optimization of control performance. Similar to that of the single-objective counterpart, metaheuristic methods are prevalently used for the multi-objective tuning of classical control, such as PID (Ayala & dos Santos Coelho, 2012; Reynoso-Meza, García-Nieto, Sanchis, & Blasco, 2013), sliding mode control (Mahmoodabadi, Taherkhorsandi, & Bagheri, 2014; Taherkhorsandi, Mahmoodabadi, Talebipour, & Castillo-Villar, 2014), as well as others (Reynoso-Meza,
A similar approach is applied in MPC tuning by using an off-the-shelf method of goal attainment (Exadaktylos & Taylor, 2010; Vega, Francisco, & Tadeo, 2008). Although more systematic than general guidelines, these methods provide non-specialized approaches that do not exploit certain characteristics of the problem and potentially require a rather exhaustive and possibly computationally impractical search to produce an optimal design set. As an alternative to approaches based on guidelines and metaheuristics, analytical methods employing problem simplifications have been proposed (Bagheri & Khaki-Sedigh, 2014; Shah & Engell, 2011; Shridhar & Cooper, 1998). However, these typically overlook some aspects of the original problem such as explicit constraint handling.

The studies discussed so far consider control performance as the sole design objective, whether with a single- or multi-objective tuning outlook. The approach separates algorithm and implementation design, revealing only half the insight in control tuning. Implementation design largely determines the hardware cost for the controller and is often not known a priori, thus is an important part of the process. A co-design approach provides a more comprehensive approach that optimizes control performance, as well as implementation cost that is dictated by the required computational capacity to implement the control in real-time. Rather than treating the required capacity as a fixed constraint, it should be co-optimized alongside control performance, avoiding system over-design or the need to re-design the system. Furthermore, previous studies (e.g. Bagheri & Khaki-Sedigh, 2014; Exadaktylos & Taylor, 2010; Júnior et al., 2014) have typically assumed the structural parameters of the controller – such as sampling rate and prediction horizon – are fixed. Nonetheless, structural MPC parameters have been shown to have an underlying role for MPC design improvement (Bachtiar, Kerrigan, Moase, & Manzie, 2015, 2016).

In light of the above discussion, the value of a co-design approach in streamlining the design process of control systems has been noted (Allison & Herber, 2014). Further, the fundamental concept of a software and hardware co-design approach for real-time optimization has been studied (Kerrigan, 2014), although analytical results to support applications in MPC are still yet to be fully developed. The main contribution of this paper is a systematic development of the optimal MPC design with a multi-objective approach. Theoretical results concerning the nature of the design problem are presented to establish certain assumptions and guarantees. These results are then used to understand the nature of the optimization problem at hand and subsequently provide conditions that a selected optimizer must satisfy in order to effectively and efficiently compute the optimal (Pareto) frontier. The approach allows the practitioner to understand the trade-off between performance and resources in structurally tuning an MPC controller for a given real-world control problem.

The paper is outlined as follows; Section 2 contains the MPC formulation studied. The proposed multi-objective MPC design approach is then presented in Section 3. Section 4 identifies the key properties of the multi-objective problem, including smoothness properties and parameter bounds. In Section 5, conditions for an effective and efficient optimizer are presented and a compliant algorithm is proposed. Section 6 considers two real-world examples to demonstrate the design approach and importance of the conditions developed for an effective optimizer. Section 7 presents conclusions of the study and potential future work.

1.1. Notational conventions and definitions

|v|_W^2 = v^T M v. ⊙ and ⊙ denotes element-wise multiplication and division, respectively. U[a, b] is a random number uniformly distributed between a and b. Unless stated otherwise, an ordered list (column vector) is boldfaced e.g. v with the size w and elements v( 1 , ..., v(w)). A set containing several ordered lists is defined in calligraphy e.g. V and V = {v_1, ..., v_N}.

2. Controller design

Consider a nonlinear dynamic plant model

\[ \dot{x} = f(x, u) \]

with states x(t) ∈ R^nx and inputs u(t) ∈ R^nu which satisfy standard properties as described in the following.

**Assumption 1.** (x, u) ↦ f(x, u) is continuous in (x, u) and globally Lipschitz continuous in x uniformly in u.

**Assumption 2.** (x, u) ↦ f(x, u) is differentiable with respect to u for all x ∈ R^nx.

Discretization is used for the purpose of digital control, such that the plant is controlled in a sampled-data fashion at sampling instants \( t_i = ih \) for \( i ∈ \mathbb{N}_{≥0} \) with sampling period h. The control command sequence is restricted to a zero-order-hold

\[ u(t) = u_i, \quad ∀ t ∈ [ih, ih + h), i ∈ \mathbb{N}_{≥0}. \]

The aim is to control the plant by applying a control law κ to regulate the model to the origin. The control law depends on the current state \( x_i := x(t_i) \) and the control design parameters \( p_i \).

Let \( p = (p_{i_1}, ..., p_{i_N}) \) contain the design parameters \( p_{i_1}, ..., p_{i_N} \) to be tuned.

In this paper, the control command is obtained by solving a finite-horizon, optimal control problem (OCP) at each sampling instant \( t_i \).

\[ (x^{∗}(·), u^{∗}(·)) = \arg \min_{x(t), u(t)} \left[ J(x(t), u(t), p) \right] \quad \text{s. t.} \quad x(0) = x_i \]

\[ \dot{x}(t) = A x(t) + B u(t), \quad ∀ t ∈ [0, T] \]

\[ x(t) ∈ [x_0, x_f], u(t) ∈ [u_0, u_f], \quad ∀ t ∈ [0, T] \]

\[ u(t) = u(kh), \quad ∀ k ∈ \mathbb{N}_{≥0}, \quad ∀ t ∈ [kh, kh + h). \]

For succinctness, the dependence of \( t ↦ x^{∗}(t) \) and \( t ↦ u^{∗}(t) \) on \( (x_i, p) \) is omitted. Consequently, \( x_i, p \) can be omitted.

The predicted states x used internally in the OCP is distinct to the actual (measured/observed) variable x, although sized equally such that \( x(t) ∈ R^nx \) and inputs \( u(t) ∈ R^nu \). Further, also note the distinction between the nonlinear plant model \( f \) and prediction model in linear time-invariant (LTI) form

\[ f = A x + B u = d_k^{\text{LTI}} \begin{bmatrix} x + d_k^{\text{LTI}} & u \end{bmatrix} \]

used internally in the OCP. The two models have the same equilibrium at the origin, that is \( f(0, 0) = 0 \). The optimization is subject to the prediction model (1c) representing the dynamics of the plant initialized at (1b), and the plant constraints (1d). The zero-order-hold control (1e) discretizes the control command over the sampling steps \( k ∈ \{0, ..., N − 1\} \).
One common choice for the OCP cost function in (1) is a quadratic

$$j(x, u, p) = \int_0^{nh} \|\dot{x}(r)\|^2_2 + \|u(r)\|^2_2 dr + \|\bar{X}(T)\|^2_2 \tag{3}$$

to penalize the state/input deviations from zero. This cost is composed by the \textit{stage cost} weighted by a positive semidefinite $Q$ and positive definite $R$, and the \textit{terminal cost} weighted by a positive semidefinite $Q_f$. $Nh := N h$ is the prediction horizon length, with $N$ prediction steps. Lastly, the OCP is to be non-degenerate.

**Assumption 3** (Non-degeneracy). The OCP (1) is non-degenerate so that its solution $(x^*, \bar{u}^*)$ is unique.

### 2.1. Design parameters

From the OCP formulation, a key design parameter is the sampling time $h$ that dictates how often a new control input can be commanded to the plant. This also sets an upper-bound on the time available for the computing hardware to solve the OCP. Next, along with the sampling time, the number of prediction steps governs the length of the prediction horizon $T := Nh$. This is the time horizon in which constraints can be applied in the prediction of the future plant behavior, and thus its value affects the performance of the controller. Further, the number of prediction steps directly affects the size of the OCP problem, that is the number of unknowns in the problem. Each of $h$ and $T$ affect both design objectives.

Generally, the prediction model type is a design parameter that can be chosen, for example, as a linear-time invariant (LTI), linear time-varying, or nonlinear model. In this study, the prediction model (1c) is defined to be LTI, a particularly application-relevant choice that reduces the general nonlinear OCP into a quadratic program (QP) for which many practical solvers exist. Furthermore, this study focuses on the MPC algorithm parameters, so that those that exclusively are attributes of the implementation, such as data precision (Kerrigan, 2014), are not considered even if they affect both objectives.

Cost function attributes, namely the cost weighting matrices, are also design parameters. These affect the OCP solution and hence the control performance, and also the time taken to numerically solve the OCP particularly when certain solver algorithms are used. For solvers which are insensitive to ill-conditioning (e.g. interior point method, Dang, Tran, & Ling, 2015), the parameters can be assumed to only affect control performance. On the other hand, the algorithm used to solve the OCP is a design choice that only affects computational complexity. As long as the solver is convergent, it is assumed that it will find the one local (thus global) minimum of the OCP as a convex problem (QP). A related parameter is the representation of the OCP, e.g. dense and sparse representations, for which some are suited to a particular algorithm and some, another. This also can be assumed to produce the same OCP solution and not affect control performance. Finally, a numerical tolerance can be used as an algorithm attribute, dictating the accuracy of the numerical solution of the OCP. This tolerance value would affect both control performance and required computational capacity.

In multi-objective design, focus shall be given to parameters that are \textit{coupled} i.e. those that affect both design objectives. In the proposed design problem, the coupled parameters of the MPC architecture, namely the sampling time $h$ and number of prediction steps $N$, are the considered tuning parameters. Conversely, parameters that only affect one of the objectives, hence \textit{decoupled}, are fixed. For the fixed parameters, in particular the solver algorithm, tolerance and numerical precision, it is assumed that they are well-chosen so that the true, global solution of the OCP can be obtained.

**Assumption 4.** The numerical solution of the OCP obtained is close to the true/analytical solution.

### 2.2. Design objectives

#### 2.2.1. Control performance

The control performance measures how well the controller steer the states to the origin over time, indicating, for example, the speed of response to set point changes, multi-variable decoupling and damping performances. This can be obtained through a simulated environment of the controlled plant. There are many measures that can be used to indicate control performance, the most common being the integrated squared error (ISE) of the states and inputs,

$$U(x_0, p) = \int_0^m v(x(r), u(i)) dr \tag{4a}$$

s. t. $x(0) = x_0 \tag{4b}$

$$\dot{x} = f(x, u) \tag{4c}$$

$$u(r) = k^s(x(x(r), p)), \quad \forall r \in [ih, ih + h), \quad \forall i \in \mathbb{N}_{>0} \tag{4d}$$

where the cost function is chosen as

$$v(x, u) = \|x(r)\|^2_2 + \|u(r)\|^2_2$$

in accordance with the OCP cost function (4). $U(x_0, p)$ defines the closed-loop value function for a given initial condition and design parameter choice.

The feasibility and stability of the closed-loop system shall be introduced as follow.

**Assumption 5** (Recursive feasibility). The closed-loop system (4) is recursively feasible such that for a given initial state $x_0 \in X_0$, the OCP (1) is feasible and remains feasible at all subsequent sampling steps. $X_0$ is the associated feasible set.

**Assumption 6** (Stability). $x_0 \in X_0$, OCP (1) guarantees asymptotic stability of the closed-loop system (4).

The assumptions allow the simulated closed-loop system to run indefinitely and will not encounter an abrupt termination (e.g. due to infeasibility of the OCP). Consequently, the ISE measure (4a) is well-defined.

**Remarks 7.** Not all choices of the parameters $(h, N)$ would validate Assumptions 5 and 6. The assumptions serve as a starting guideline when choosing the parameters for the closed-loop system to be functional (feasible and stable). The set containing such parameters is to be narrowed further to define those that are optimal.

Specifying the exact conditions to guarantee Assumptions 5 and 6 is not the focus of this study and the reader shall refer to the relevant references that are amply available. A brief guideline is given below.

1. Levis, Schlueter, and Athans (1971) note that there exist critical sampling periods where the discrete plant loses full controllability, which consequently leads to an unstable closed-loop trajectory. Since fast systems are of interest, short sampling periods can be appropriately considered (shorter than the first critical sampling periods).
2. For $f = f$, an appropriate choice of $Q_i$ (Mayne, 2014) and $T$ (Kerrigan & Maciejowski, 2000) is able to guarantee recursive feasibility and stability. The restriction on model fidelity might be relaxed so that the discrepancy between the two is bounded, instead of zero.

If either one of the Assumptions 5 and 6 is not valid, a penalised value of the metric $U$ can be set for cases when the simulation terminates early.

A linear combination of the closed-loop value function

$$V(X_p, p) = \sum_{i=1}^{T} w_i U(X_0, p) \quad \forall X_0 \in X_0$$

(5)

provides a numerical measure of the control performance of the MPC law (2) with the OCP (1). A linear combination is chosen to preserve the smoothness properties of the value-function $U$. $X_0$ is a set of initial conditions representative of the intended operating range of the controlled plant, each of which is weighted by $w_i$ to determine the relative significance of each scenario.

### 2.2.2. Required computational capacity

In designing the controller, the capability of the computational hardware onto which the controller will be implemented is often not known a priori. It is therefore imperative to set the required computational capacity, which is a primary factor in determining cost, as a design objective. Structural MPC parameters dictate the computational complexity of the OCP (1). In turn, the complexity governs the required resource, affecting the time taken by a processing unit to generate a control command that is upper-bounded by the sampling period.

In many cases, a numerical simulation is conducted instead of directly testing the control plant to test the performance of the control architecture. The computational data obtained in simulation is reflective of the hardware used in the simulation platform, namely the simulation hardware. This is to be differentiated from the implementation hardware, which is the actual hardware used to implement the controller and control the plant.

Let the upper-bound for the time taken for a command input to be generated using the simulation hardware be $\gamma(p)$. This indicates the computational complexity for a given design $p$. Based on this, a dimensionless measure denoted the Capacity Number can be derived as

$$\eta(p) = \frac{\gamma(p)}{h}$$

(7)

to indicate the required power of the implementation hardware relative to the simulation hardware so that the MPC controller can be functionally implemented.

For the Capacity Number to be a meaningful measure, an assumption which specifies the relationship between the simulation and implementation hardware is needed.

#### Assumption 8 (Scalability). The solution time upper-bound of the simulation hardware $\eta$ is linearly scalable to the solution-time upper-bound $\eta_i$ if the implementation hardware was used. That is, $\eta = a\eta_i$ for some constant multiplier $a \in R_{>0}$.

The importance of the scalability Assumption 8 will be clarified later (Remark 11) when the multi-objective design problem is formulated.

The solution time upper-bound $\gamma$ depends on the parameters of the OCP as a QP. For the positive definite cost function (3), QP optimization is known to be a P problem solvable in polynomial time as its size increases (Kozlov, Tarasov, & Khachiyan, 1980). The size of the OCP is directly associated with the number of unknowns, which is directly proportional to the number of prediction steps $N$. Consequently, the solution time of the OCP as a quadratic program will be modelled as a polynomial with respect to $N$.

#### Assumption 9 (Solution time upper-bound). The upper-bound on solution time is monotonically increasing with the number of prediction steps $N$, modeled by a polynomial of degree $n$

$$\gamma(p) = \sum_{i=0}^{n} a_i N^i$$

(6)

for some constants $a_i, i \in \{0, \ldots, n\}$.

Generally, the parameters $a_i$ are dependent on design parameters such as sampling time, OCP representation and solver algorithm. For instance, sampling time has an effect on the matrices of the OCP and how well they are conditioned, thus affecting solution time. However, variations in solution time from difference in sampling time are not significant in the tests conducted in this study and are treated as negligible. Furthermore, since other design parameters are treated as constants, the parameters $a_i$ are constant.

Assumption 9 is later verified in the results in Section 6. Consequently,

$$\eta(p) = \frac{\gamma(p)}{h} = \sum_{i=0}^{n} a_i N_i$$

(7)

### 3. Multi-objective design of MPC

Let the objectives be defined as a vector $\varepsilon=(V, \eta) \in \Re^{n_e}$ where $n_e = 2$. The multi-objective design (MOD) of MPC is posed as the following:

$$\underline{\varepsilon}^* (\mathcal{P}_1) = \arg \min_{p} \{ \eta(p) \}$$

(8a)

s. t. $p \in \mathcal{P}_1$

(8b)

The minimization (denoted m-min) is a multi-objective minimization to find the Pareto optimal design set $\mathcal{P}_1$ for a given search space $\mathcal{P}_1$ (see Remark 7). This solution set contains the Pareto optimal design choices for a practitioner to select from, based on the Pareto front

$$\mathcal{L} (\mathcal{P}_1) = \{ (\varepsilon(p) | p \in \mathcal{P}_1) \}$$

(9)

that shows the optimal trade-off between the two objectives. The Pareto optimal (non-dominated) points, defined below, make up the Pareto front.

#### Definition 10 (Pareto optimal point (Deb, 2001)). A point $p_1 \in \mathcal{P}$ is a Pareto point if there does not exist another design choice $p_2 \in \mathcal{P}$ such that $p_2$ dominates $p_1$.

The optimization is contrasted to single-objective optimization where the solution correspond to one point (the minimum).

#### Remarks 11. Assumption 8 implies that a Pareto optimal point obtained using the simulation hardware stays Pareto optimal for the implementation hardware.

In the proposed tuning problem, focus is given on the two underlying structural design parameters of the OCP, sampling time and prediction horizon. That is,

$$p=(h, N) \in \mathcal{P} := \Re_{>0} \times \Re_{>0}$$

(10)
The rest of the coupled design parameters and all the decoupled design parameters are fixed. With such a focus, the resulting design problem would have specific characteristics, as revealed next, allowing for a specialized optimizer algorithm to be proposed.

4. Key properties of the MOD-MPC system

This section analyzes the MOD-MPC system to reveal its key attributes. These include the smoothness properties and bounds on the objective and design parameters, establishing the underlying assumptions and guarantees that are useful for the subsequent development of a numerical optimizer. Based on the analytical foundation, a targeted solution method that is both effective and efficient can be appropriately developed.

4.1. Smoothness of the design objectives

4.1.1. Monotonicity

The model used for the required computational capacity $\eta(\cdot)$ has the following monotonicity property.

**Proposition 12 (Monotonicity of $\eta$).** Consider the required computational capacity $\eta(\cdot)$ in (7) and that Assumption 9 holds. For $\mathbf{p} = (h, N)$, $\eta(\mathbf{p})$ is:

- monotonically decreasing with respect to $h$,
- monotonically increasing with respect to $N$.

**Proof.** The monotonicity of $\eta$ can be directly taken from the dependence of $\eta$ to $h$ and $N$ as given in (7).

On the other hand, the control performance measured by the closed-loop value function $V$ is non-monotonic. The non-monotonicity, defined next in Propositions 13 and 14, is demonstrated numerically in Bachtiar et al. (2015, 2016).

**Proposition 14 (Non-monotonicity of $V$ with respect to $h$).** Consider the control performance $V(x_h, \cdot)$, $x_h \subseteq x_f$ as in (5) and that Assumptions 3–6 hold. For $\mathbf{p} = (h, N)$, $h \rightarrow V(x_h, \mathbf{p})$ is generally non-monotonic. Specifically, the sets $H$, and $H$, such that $h \rightarrow V(x_h, \mathbf{p})$ is monotonically increasing $\forall h \in H$, and monotonically decreasing $\forall h \in H$.

**Proof.** Consider the closed-loop system. For a fixed $N$, increasing $h$ both desirably increases the prediction horizon length $T = Nh$ but also undesirably slows down the sampling rate $1/h$ of the closed-loop system. At the limit $h \rightarrow 0$, the system operates badly, since the control prediction barely, if at all, captures any dynamics of the system. As $h$ is increased, performance is improved, until some point where all important dynamics are captured. Increasing prediction length is no longer as influential as the delayed sampling rate, after which the performance is worsened with increasing $h$. Therefore, the value function $U(x_h, \cdot, N)$ is generally non-monotonic with respect to $h$. $V$ is a linear combination of $U$ and thus obeys the same properties. Consequently, there are two mutually exclusive sets $H$, and $H$, such that for all $h_a \in H$, and $h_b \in H$, if $h_b \geq h_a$ then $V(x_h, (h_a, N)) \geq V(x_h, (h_b, N))$. $H$ is defined similarly.

**Proposition 13 (Non-monotonicity of $V$ with respect to $N$).** Consider the control performance $V(x_h, \cdot)$, $x_h \subseteq x_f$ as in (5) and that Assumptions 3–6 hold. For $\mathbf{p} = (h, N)$, $N \rightarrow V(x_h, \mathbf{p})$ is generally non-monotonic. Specifically, the sets $N_h$, and $N_f$, such that $N \rightarrow V(x_h, \mathbf{p})$ is monotonically increasing $\forall N \in N_h$, and monotonically decreasing $\forall N \in N_f$.

**Proof.** Consider the closed-loop system (see footnote 1). For a fixed $h$, increasing $N$ increases the prediction horizon length $T = Nh$. As a result more dynamics, as well as plant-model mismatch, are captured by the prediction. There is a trade-off balance between the two so that the value function $U(x_h, (h, \cdot))$ is generally non-monotonic with respect to $N$. $V$ is a linear combination of $U$ and thus obeys the same properties. Consequently, there are two mutually exclusive sets $N_h$, and $N_f$, such that for all $N_a \in N_h$, and $N_b \in N_f$, if $N_b \geq N_a$ then $V(x_h, (h, N_a)) \leq V(x_h, (h, N_b))$. $N_f$ is defined similarly.

4.1.2. Continuity and differentiability

The continuity of the solution of the OCP (1) with respect to $\mathbf{p}$ is described in the following.

**Lemma 15 (Continuity of $u^*$ (Bachtiar et al., 2015, 2016)).** Consider the OCP (1) satisfying Assumptions 3–6. The unique optimal solution $u^*(x_h, \mathbf{p})$ of the OCP is differentiable with respect to $h$ for a given $N$.

**Proof.** The proof is given in Lemma 14 in Bachtiar et al. (2016).

**Theorem 16 (Continuity of $V$).** Consider the control performance $V(x_h, \mathbf{p})$, $x_h \subseteq x_f$, and that Assumptions 3–6 hold. If Assumptions 1 and 2 hold, $h \rightarrow V(\mathbf{p})$ is continuous with respect to $h$.

**Proof.** Let $f(\cdot, \cdot)$ satisfy Assumption 1. Let $z(\cdot, \cdot)$ be the solution of $x = f(x, u)$, $u = z(t, u)$ is continuous with respect to $u$ (Theorem 3.5 in Khalil, 2002). The control law (4d) is given by $u(t) = u(t) = u(t) = u(t)$, $\forall e \in [\tilde{t}, h + h]$, $\forall v \in N_{x_b}$. Subsequently, from Lemma 15, $u$ is continuous and differentiable with respect to $h$, implying that the solution $h \rightarrow z(t, u(x_h, (h, N)))$ is continuous with respect to $h$ for a given $N$. $h \rightarrow U(x_h, (h, N))$, $\forall x_h \in x_h$ is thus continuous with respect to $h$ for a given $N$ and satisfaction of Assumptions 3–6. $V$ is a linear combination of $U$ and thus has the same monotonicity properties.

The differentiability of the closed-loop value function can be described and is stated in the following.

**Theorem 17 (Differentiability of $V$).** Consider the control performance $V(x_h, \mathbf{p})$, $x_h \subseteq x_f$, and that Assumptions 3–6 hold. If Assumptions 1 and 2 hold, $p \rightarrow V(\mathbf{p})$ is differentiable with respect to $h$ for a given $N$.

**Proof.** Let $f(\cdot, \cdot)$ satisfy Assumptions 1 and 2. Let $z(\cdot, \cdot)$ be the solution of $x = f(x, u)$, $u = z(t, u)$ is differentiable with respect to $u$ (Theorem 3.5 and Section 3.3 in Khalil (2002)). The control law (4d) is given by $u(t) = u(t) = u(t) = u(t)$, $\forall e \in [\tilde{t}, h + h]$, $\forall v \in N_{x_b}$. Subsequently, from Lemma 15, $u$ is continuous and differentiable with respect to $h$, implying that the solution $h \rightarrow z(t, u(x_h, (h, N)))$ is differentiable with respect to $h$ for a given $N$. $h \rightarrow U(x_h, (h, N))$, $\forall x_h \in x_h$ is thus differentiable with respect to $h$ for a given $N$ and satisfaction of Assumptions 3–6. $V$ is a linear combination of $U$ and thus has the same monotonicity properties.

4.2. Competing nature of the design objectives

The competing nature of a pair of functions that are both to be minimized (or maximized) is defined as follows.

**Definition 18 (Competing functions).** Two functions $a \rightarrow f(a)$ and $a \rightarrow g(a)$ are competing with each other for the design set $[g, \pi]$ if $a \rightarrow f(a)$ is monotonically increasing and $a \rightarrow g(a)$ is monotonically decreasing (or vice versa) on $[g, \pi]$.

In the multi-objective design of MPC, both the closed-loop value function $V$ and required computational capacity $\eta$ are to be
minimized. Based on Propositions 12–14, the two design objectives of control performance and required computational capacity are competing. This is detailed in the following.

**Lemma 19 (Competing design objectives).** The objective functions \( V(x_0, \cdot), x_0 \in X_0, \) and \( \eta(\cdot) \) are competing as per Definition 18 within the design parameter set \( \mathcal{P} = (h, x) \) from Proposition 12–14, and given that Assumptions 3–6 and 9 hold.

**Proof.** Consider Proposition 12–14 and satisfaction of Assumptions 3–6 and 9. \( \eta \) is monotonically decreasing with respect to \( h \) and increasing with \( N \). \( V \) is non-monotonic with respect to \( h \) and \( N \), and there exist a set such that \( V \) is increasing with respect to \( h \) and decreasing with respect to \( N \), given by \( \mathcal{P} = (h, x) \) from Propositions 13 and 14. Within this set, \( \forall p \in \mathcal{P}, \ p \to V(p) \) is monotonically increasing whilst \( p \to \eta(p) \) is monotonically decreasing, or vice versa.

The search space \( \mathcal{P} \) is assumed to intersect with \( \mathcal{P}_c \) so that the MOD-MPC solution \( \mathcal{P}_s \) exist.

**Assumption 20.** \( \mathcal{P}_c = \mathcal{P} \cap \mathcal{P} \neq \emptyset \).

The associated Pareto front \( \mathcal{L} \) as per (9) consists of Pareto optimal points as defined in Definition 10, each of which is a Pareto design choice \( p \in \mathcal{P} \subseteq \mathcal{P}_c \).

Further to its Pareto optimality, the quality of a point can be specified by its rank, as defined in the following.

**Definition 21 (Rank).** Given a countable set of points \( \mathcal{P} \), if a point \( j \) is Pareto optimal then its rank \( r_1 \). Subsequently, a point \( j \) has rank \( r_j \) if it is Pareto optimal in \( \mathcal{P} \setminus \mathcal{P}_{j-1} \), where \( \mathcal{P}_j \) is the set of all points with rank \( r \leq r_j \).

That is, all Pareto optimal points in a given set of points have a rank of 1. The Pareto optimal points in the set that excludes points with rank 1 have rank 2, and so on.

### 4.3. Bounds on the pareto design set

The first bounding of the design parameter comes from the fact that it is numerically impractical to search the open set \( \mathcal{P} \) in (10). Hence, the search space \( \mathcal{P}_c \) in (8) must be a closed set \( \mathcal{P}_c \subseteq \mathcal{P} \) that is able to be practically searched to find the Pareto design set \( \mathcal{P}_s \).

Next, an assumption on the Pareto design set can be made based on some intuitions on the nature of the design problem.

**Assumption 22 (Upper-bound on \( h \)).** For a given \( N > 1 \), the Pareto design set is upper-bounded by \( \bar{h} \). This bound is defined by the notion that \( \exists \bar{h} \in \mathcal{P}_c \) such that \( \forall h > \bar{h}, \ \exists \bar{h} < h \) such that \( \mathcal{E}(h, N - 1) < \mathcal{E}(h, N) \).

The assumption comes from the fact that as the sampling period \( h \) is increased for a given number of prediction steps \( N \), the competitive effect of reducing the required computing resource will be diminished and surpassed by the deterioration in control performance as a result of more infrequent sampling. At this point, using fewer prediction steps, e.g., \( N - 1 \), would reduce resource requirement more competitively than deteriorating performance. The opposite is also true; as the sampling period is shortened, the effect of increasing control performance due to a finer sampling would eventually be overtaken by an increase in computing resource required, after which increasing \( N \), to e.g., \( N + 1 \), would be more competitive in improving performance.

**Assumption 23 (Lower-bound on \( h \)).** For a given \( N \geq 1 \), the Pareto design set is lower-bounded by \( \hat{h} \). This bound is defined by the notion that \( \exists \hat{h} \in \mathcal{P}_c \) such that \( \forall h < \hat{h}, \ \exists \hat{h} < h \) such that \( \mathcal{E}(h', N + 1) < \mathcal{E}(h, N) \).

Based on the bounds on the sampling period and the smoothness properties of the design objectives defined in the earlier subsection, bounds can be specified for \( \mathcal{P}_s \).

**Theorem 24 (Bound on \( \mathcal{P}_s \).** Consider a rectangular search space \( \mathcal{P} = (h, x) \) and that Assumptions 22 and 23 is satisfied. The Pareto design set \( \mathcal{P}_s \) is contained within the band \( \mathcal{P}_c = \{ p : h < m_0 N + h_0 \text{ and } h < m_0 N + h_0 \} \) for some negative gradients \( m_0 \in \mathbb{R}^c_\alpha \) and constants \( h_0, i \in \{1, 2\} \).

**Proof.** Consider a Pareto design set \( \mathcal{P}_s \) in a rectangular search space \( \mathcal{P} = (h, N) : h \in [\bar{h}, \bar{h}] \text{ and } N \in [N_i, N] \text{ and } N \). For a given number of prediction steps \( N > 1 \), let the sampling periods corresponding to the Pareto design set have an upper-bound from Assumption 22 denoted \( \bar{h}_N \). Paraphrasing the assumption, \( \exists h \text{ such that } \mathcal{E}(h, N - 1) < \mathcal{E}(h, N), \forall h > \bar{h}_N \) for the given upper-bound \( \bar{h}_N \). Consequently, there must exist an upper-bound associated with \( N - 1 \) prediction steps \( \bar{h}_{N-1} \) that is larger than the upper-bound \( \bar{h}_N \), giving \( \bar{h}_N < \bar{h}_{N-1} \). Therefore, the Pareto design set can be upper-bounded by a line of negative gradient with respect to \( N \). An opposite notion can be made using Assumption 23 to form a lower-bound with a negative gradient. This gives a bound in the form of a band \( \mathcal{P}_s \) as in the theorem.

### 5. Numerical solution to the MOD-MPC problem

#### 5.1. Conditions for an effective and efficient optimizer

Analysis of the key properties of the system results in a number of characteristics required by a proposed numerical optimizer used to solve the MOD-MPC problem (8) accurately and quickly, as summarized below.

**Condition 1 (Convergent).** Lemma 19 implies that there is a Pareto optimal design set \( \mathcal{P}_s \) for a given search space \( \mathcal{P} \), associated with the trade-off of the competing objectives. The optimizer should be able to effectively find \( \mathcal{P}_s \) with certain guarantees.

**Condition 1a (Global).** Propositions 12–14 define that the required computational capacity \( \eta \) is monotonic and that the value function \( V \) is non-monotonic. The optimizer needs to search globally and handle the many local optima on the objective surface.

**Condition 1b (Able to handle discrete parameters).** The optimizer must be able to handle discrete design parameters defined in (10).

**Remarks 25.** Conditions 1a and 1b are necessary conditions for Condition 1 to be fulfilled. Satisfaction of these two is not always sufficient to satisfy Condition 1.

Condition 1 (necessarily with Conditions 1a and 1b) is a sufficient condition for a numerical optimizer to be accurate (convergent) for the MOD-MPC problem. Additional features are necessary for the optimizer to converge quickly and efficiently.

**Condition 2 (Continuous).** The optimizer could rely on Theorems 16 and 17 that define the continuity/differentiability of the value function \( V \) based on the knowledge of \( \mathcal{P}_s \).

**Condition 3 (Focused).** Theorem 24 states that for a rectangular search space, the Pareto optimal design set is located within a specific space defined as a band. As a consequence, the optimizer should be able to focus its search within the band and omit any ineffectual space.

Conditions 2 and 3 are sufficient for optimizer efficiency and
better performance than a general-purpose method.

5.2. A compliant optimizer algorithm (DITRI)

Having now specified the sufficient conditions for effectiveness and efficiency, a specialized optimizer satisfying all the conditions can be proposed for the MOD-MPC problem. An algorithm is proposed based on Lipschitzian optimization (Jones, Perttunen, & Stuckman, 1993), denoted ‘Dividing TRIangles’ (DITRI), with details outlined in Appendix A.

The choice of a Lipschitzian approach addresses Conditions 1a and 1b necessary for the optimizer to be convergent. Lipschitzian optimization is gradient-free and is built for a global search, therefore addressing Condition 1a. The method directly handles discrete design parameters specified in Condition 1b. As a whole, the proposed optimizer is guaranteed to converge and satisfies Condition 1, making it an accurate optimizer that can effectively find the solution (Pareto design set) \( P^* \) of the MOD-MPC problem (8) in a given search space.

**Theorem 26** (Convergent DITRI). Consider a multi-objective optimal design problem with objectives \( p \rightarrow e(p) \) for \( p \in P_1 \) where \( P_1 \) is a finite search space. Let the solution of the problem be \( P_1 \subseteq P_2 \). Also let \( S_i \) for some indexing variable \( i \) be a sequence of solutions generated from a given initialization \( S_0 \), DITRI, described in Algorithm 2, is a convergent algorithm such that \( \lim_{i \to \infty} S_i = P_1 \) for some \( i \to \infty \).

**Proof.** The proof is given in Appendix B.\(^\square\)

As well as being accurate, DITRI is designed to be an efficient optimizer by using projection of bounds (A.1). This assumes continuity in the objective function and satisfies Condition 2. Finally, DITRI conducts a focused search of a given search space as outlined in Appendices A.2–A.4 to fulfill Condition 3.

6. Simulation results

Two real-world examples are investigated. \( Q \) and \( R \) are chosen accordingly for each case. \( Q = P_{\text{MAX}} \) where \( P_{\text{MAX}} \) is the solution of the algebraic Riccati equation for the simulated plant. The OCP is represented as a sparse QP. The global search criterion (A.3) is set as \( \tilde{d}(i) = \frac{5}{2\sqrt{3}} \) so that no search space is larger than that equivalent to \( i^8 \) divisions from the initial triangle.

The simulation hardware specifications are given in Table 1. Gurobi (Gurobi Optimization, 2013) in MATLAB is used in this study as the OCP solver. The algorithm used is the interior point (barrier) method, with all tolerances set as the default. A representative result for the relationship between the solution time and number of prediction steps \( N \) for a range of sampling time \( h \) is shown in Fig. 1. It is shown that \( \gamma \) is generally increasing with \( N \) and that it is very weakly correlated to \( h \), verifying Assumption 9. From the obtained data, the relationship for the chosen algorithm and QP form is mostly linear. \( \gamma \) in (6) is modeled as a linear function with the chosen constants \( a_0 = 1.3 \times 10^{-4} \) and \( a_0 = 3.5 \times 10^{-3} \). The relevant data and model is shown on the top-left graph of the figure.

### Table 1

<p>| | | |</p>
<table>
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<td><strong>FLOP/s available for simulation</strong></td>
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</tbody>
</table>

6.1. Test plant models

The first test case looks at a missile pitch-axis autopilot (PAA) adapted from Bachtiar et al. (2014). The missile is flown at a cruising altitude and the autopilot is to control the missile to track a commanded acceleration. The second case aims to tune a controller for diesel engine control. The engine is modeled by a mean-value engine model (MVEM) taken from Broomhead, Manzie, Brear, and Hield (2015) and the control objective is to track a given engine speed and power output.

A coordinate shift is applied appropriately to transform a given tracking problem into that of regulation,

\[
x = e - e_s \quad x_o = e_0 - e_s \quad \mathbf{x} = e - e_s (x + e_o) = f(x, u) \quad \text{etc. (11)}
\]

6.1.1. Pitch-axis autopilot

The first test case looks at a missile pitch-axis autopilot (PAA) at 20 000 ft. The nonlinear tracking model is

\[
f(e, u) = \begin{bmatrix}
e_2 + \cos(e_1) f(e, u)/(\text{in}) \\
L(e, u)/(f_y') \\
e_4 \\
-2a_0e_3 - 2a_0e_4 + a_0 e_5 \\
u
\end{bmatrix}
\]

(12)

where \( F \) and \( L \) are the nonlinear mapping for the aerodynamic lift force and pitching moment respectively. \( e_1 \) is the angle of attack and \( e_2 \) is the pitch rate of the missile. The actuation of the fin deflection \( e_2 \) is modeled as a second order system. The input \( u \) is the rate of the commanded fin deflection \( e_2 \). Missile speed \( v = Mv_\gamma \), where \( v_\gamma \) is the speed of sound at 20 000 ft, is constant at Mach number \( M = 0.25 \). \( m \) and \( I_y \) are the mass and moment of inertia of the missile respectively. These parameter values, other missile frame parameters, constants related to the actuation dynamics, along with the aerodynamic coefficients and models used for \( F \) and \( L \) are the same as given in Bachtiar et al. (2014); Nichols et al. (1993).

The control objective is to track a given acceleration output \( y = F/(mg) \), where \( g \) is the gravitational acceleration. The test scenario is to track 5 different acceleration outputs from steady-state at \( 0 \), \( e_o = 0 \). The outputs are \( 2, 4, 6, 8 \) and \( 10 \), each associated with a unique steady state \( e_s \) and initial condition as per (11), making up the set of initial conditions \( X_0 \) in (5) that are equally weighted, \( w_i = 0.2 \), \( i \in \{0, \ldots, 5\} \). The states are upper- and lower- bounded by \( e_s = \{20^\circ, 35^\circ/s, 45^\circ, 10^6/s, 45^\circ\} \) and \( g = -\mathbf{g} \). The input is bounded by \( u = 10^6/s \) and \( y = -\pi \). \( Q = C^TC \) where \( C = dy/\Delta t \mathbf{e}_{\text{yo}} \) comes from the linearization of the output \( y \) at the target steady state and \( R = 10^{-6} \).

6.1.2. Diesel engine control with a mean-value engine model

The second case looks at engine control with a 5-state 3-input mean-value engine model (MVEM) from Broomhead et al. (2015). The 5 states are the engine speed \( e_3 \), turbine speed \( e_2 \), VGT actuator position \( e_1 \), intake manifold pressure \( e_4 \) and temperature \( e_5 \). The three inputs are the injection duration \( u_1 \), load applied to the engine by the generator \( u_2 \) and the VGT commanded position \( u_3 \). The model is

\[\text{Symbols used in a plant model are used exclusively in the model and should not be confused with symbols introduced elsewhere.}\]
The initial state of the engine is at 2000 rpm producing 20 kW of power. The control objective is for the engine to track a steady-state at 2500 rpm producing 36 kW of power. The states are constrained with an upper-bound of \( e = (2500 \text{ rpm}, 150 000 \text{ rpm}, 10^8, 10^8) \) and lower-bound of \( e = (1500 \text{ rpm}, 45 000 \text{ rpm}, 0, 0, 0) \). The input is bounded by \( \Pi = (1 \text{ ms}, 300 \text{ Nm}, 90) \) and \( y = (0.5 \text{ ms}, 10 \text{ Nm}, 60) \) respectively. \( Q = \text{diag}(1, 0, 0, 0, 0) \) and \( R = \text{diag}(0, 1, 1) \) as to track engine speed, power output, and VGT position.

6.2. Test results

Fig. 2 shows a representative result for the PAA case (12). A resulting trade-off curve is obtained after 20 evaluations using DITRI with \( \mathbf{p} = (0.015, 15) \) and \( \mathbf{p} = (0.001, 3) \), consisting of 10 different designs. For the MVEM case (13), a representative result is shown on Fig. 3 for 20 evaluations using DITRI with \( \mathbf{p} = (0.4, 10) \) and \( \mathbf{p} = (0.05, 1) \). After 20 evaluations, 11 designs on a trade-off curve are obtained. The associated solution in the design parameter space is shown on the bottom plots of each figure.

The trade-off curves represent the set of optimal designs a practitioner can choose from. For example, design 16 of the PAA case (Fig. 2) has a controller design with a sampling time of \( h = 6.5 \text{ ms} \). Consequently, the implementation hardware of the controller should be able to solve the OCP with \( N = 11 \) prediction steps (prediction horizon of \( \sim 70 \text{ ms} \)). Consequently, the implementation hardware of the controller should be able to solve the OCP with \( N = 11 \) in under 6.5 ms.

To help design the implementation hardware, the Capacity Number of the design choice can be examined. PAA design 16 is associated with \( \eta = 0.82 \), indicating that the implementation hardware must have at least 0.82 times the processing power in FLOP/s of the simulation hardware (Table 1). This depends on hardware capabilities and implementation architecture, including
The trade-off curves also reveal the sensitivity of control performance to computational capacity. In the PAA case, performance improvement after $\eta \approx 0.6$ is not significant anymore. This implies that there is not much benefit to be gained from hardware more powerful than $\eta \approx 0.6$. In the MVEM case, the value is $\eta \approx 0.015$.

6.3. Validation of the prescribed optimizer conditions

To show the importance of the conditions for effectiveness and efficiency prescribed in Section 5, DITRI is compared to two other algorithms. The first is a non-dominated sorting genetic algorithm (NSGA) adopted from Deb, Pratap, Agarwal, and Meyarivan (2002) and the second is one based on surrogate hyper-volume improvement (HVOL) from Tesch, Schneider, and Choset (2013). Table 2 outlines how each optimizer satisfies the specified conditions.

6.3.1. Effective convergence (Condition 1)

To assess the convergence of the trade-off curves obtained, the curves are compared to the true Pareto front $\mathcal{L}$. Since the true Pareto front is not known, it is approximated by doing a full exploration on a uniform grid of 400 $h$-values for each $N$-value in the parameter space. For both case studies, it is shown that the solution obtained by DITRI is close to the true Pareto front of the problem (Figs. 2 and 3).

The closeness of a trade-off curve to the true Pareto front can be measured by calculating the average of the closest Euclidean distance between each point on the trade-off curve to the Pareto front. This measure is denoted $\Delta$ and plotted in Fig. 4 against function evaluation count. A second metric calculates the Euclidean distance of the tips (vertices) of the trade-off curve and Pareto front. This measures the coverage of the solution, denoted $\Psi$, and is shown in Fig. 5. Calculation of both metrics are based on a normalized design objective values.

Figs. 4 and 5 show that trade-off curves produced by DITRI and NSGA approach the Pareto front with increasing function evaluation counts. However, HVOL struggles to converge. Figs. 2 and 3 show non-convergent trade-off curves, each from 100 evaluations using HVOL, confirming HVOL’s inability to find the Pareto front. This is consistent with the expectation, since both DITRI and NSGA satisfies Condition 1 for convergence, whereas HVOL does not. The fulfillment of Conditions 1a and 1b by HVOL is not sufficient to guarantee convergence, consistent with Remark 25.

6.3.2. Efficient search (Conditions 2 and 3)

Results in Figs. 2 and 3 show that the Pareto solution lies within a band defined in Theorem 24. DITRI takes advantage of this. Comparing the results in Fig. 4, DITRI exhibits the best performance.
convergence rate, consistent with the fact that DITRI satisfies Conditions 2 and 3 for efficiency, while the general-purpose NSGA and HVOL do not. In Fig. 5, the indication that HVOL is quick to reach the Pareto front tips shall be taken with full awareness as HVOL does not have convergence guarantees.

7. Conclusions and future work

This paper presented a multi-objective MPC tuning approach, treating control performance and the required computational capacity as concurrent objectives in a given control problem. Focus was given to the design of structural MPC parameters, namely sampling time and prediction horizon. The approach is an extension to those that explore only one design objective and treat algorithm/software and implementation/hardware separately. A co-design of algorithm and implementation streamlines the design process, avoiding unnecessary costs. The proposed approach was studied analytically to present several theoretical results that reveal key properties of the design problem and subsequently prescribe necessary and sufficient conditions for an effective optimizer. Finally, two tests on real-world examples were conducted to demonstrate the design approach, as well as the importance of the conditions specified for an effective optimizer of the design problem.

Future work following the study could consider other coupled design parameters currently kept constant, such as the choice of prediction model. Furthermore, the scope can be extended further beyond the MPC structure to include attributes of the numerical method used to solve the OCP, such as the algorithm and its tolerances, as well as features of the implementation hardware resource such as data representation type. This would extend the idea to a full co-design approach that looks at attributes of both software and hardware instead of focusing only on the former.

Appendix A. DITRI algorithm

A.1. Projection of bounds

The principles of Lipschitzian optimization are outlined in Algorithm 1. In each iteration, given a set of point(s) \( P \), a point \( p \in P \) is potentially optimal if its projected lower-bound of the (minimized) objective within the associated search space must be equal or better than all points in \( P \). The search space of all potentially optimal points will be partitioned into smaller divisions and a new point is evaluated in each division.

Algorithm 1. Generic Lipschitzian optimization.

\[
\begin{align*}
\text{Require:} & \quad \text{Search space bounds} \\
1: & \quad \text{Evaluate initial point(s)} \\
2: & \quad \text{repeat} \\
3: & \quad \text{Find a set } O \text{ of potentially optimal points} \\
4: & \quad \text{for all } o \in O \text{ do} \\
5: & \quad \text{Evaluate new points based on the search space division of potentially optimal point } o \text{ do} \\
6: & \quad \text{end for} \\
7: & \quad \text{until iteration or evaluation count limit is reached}
\end{align*}
\]

The bound projection must be consistent throughout, dictated by a constant referred to as the Lipschitz constant, hence the name of the algorithm. The projected bound \( L(d_j) \) of an evaluated objective \( \mathcal{O}(p_j) \) in a search space of size \( d_j \) is defined as \( L(d_j) = \mathcal{O}(p_j) - K_d d_j \) for some Lipschitz constant \( K_d > 0 \). For the projection to make sense, it assumes that the objective is continuous, as given in Condition 2. For a point \( p \), the bound is better (smaller) if it has a smaller objective \( \mathcal{O}(p) \) and/or a bigger search space size \( d_j \). The point \( j \) is potentially optimal if \( L(d_j) \leq L(d_i) \) for all \( i \in P \), that is

\[
\mathcal{O}(p_j) - K_d d_j \leq \mathcal{O}(p_i) - K_d d_i, \quad \forall i \in P \text{ for a given } K_L > 0.
\]

A.2. Potential optimality classification

Potential optimality of a design choice is classified by its Pareto
optimality Definition (10). More precisely, the classification is based on the rank $r$ of the point (Definition 21). The specification of the Lipschitz constant is tightened from $K_i > 0$ to $K_i = \varepsilon$ where $\varepsilon$ is a very small positive number, giving

$$r_j - K_d d_j \leq r_j - K d_j \forall j \in P, \ K_i = \varepsilon.$$  

(A.2)

This criterion helps to quickly localize regions of optimal solutions, ultimately allowing for a more efficient convergence given in Condition 3.

Potential optimality selection can be intuitively illustrated on an f-d plot (Fig. A1). A point $j$ satisfies (A.1) if there is a line intersecting the point with a gradient $K_j > 0$ such that all other points lie above the line in f-d coordinates. Consequently, all potentially optimal points lie on the lower right edge of the convex hull of the points. The tightened requirement to $K_i = \varepsilon$ in (A.2) effectively means that only points with the lowest objective values (rank) are chosen to be potentially optimal.

In addition to (A.2), the criterion

$$d_j > \tilde{d}(i)$$  

(A.3)

is used to guarantee convergence in Condition 1. $\tilde{d}(i)$ is monotonically decreasing with $i$ and $\lim_{i \to \infty} \tilde{d}(i) = 0$. This asserts that search spaces that are relatively much bigger are divided, effectively acting as global search. Consequently, eventually all the search space divisions will be divided regardless of satisfaction of (A.2).

A.3. Search space normalization, bounds, division and size

The search space is bounded rectangularly, specified by $p$ and $\bar{p}$ containing the lower- and upper-bound for each design parameters. A normalized point $c$ is defined as

$$c = (p - p) \otimes (p - \bar{p})$$  

so that $p = c + (p - p) \otimes c$.  

(A.4)

Theorem 24 allows for improving search efficiency by focusing on the band $\mathcal{B}_i$ defined in the theorem. To efficiently locate the specified band, the search will be simplex based (triangular) as illustrated in Fig. A2. In each iteration, every potentially optimal simplex is divided to form two simplexes of equal size. The size measure $d$ used in criteria (A.2) and (A.3) is the longest distance from the center to the vertices of the simplex.

<table>
<thead>
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<th>Scenario 1</th>
<th>Scenario 2</th>
<th>Scenario 3</th>
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<td><img src="image2.png" alt="Diagram" /></td>
<td><img src="image3.png" alt="Diagram" /></td>
</tr>
</tbody>
</table>

Fig. A3. Comparison between the more adaptable and effective simplex based search compared to rectangular search. Each evaluation point is placed in the middle of the search space. $\mathcal{P}_b$ is defined in Theorem 24.

The efficiency of a simplex based search comes from the fact a simplex is the basic polytope in any n-dimension. The approach is contrasted to the classical implementation of Lipschitzian optimization whereby hyper-rectangles are used (DIRECT Jones, 2001). The efficiency of DITRI is demonstrated in Fig. A3. DITRI does the minimum evaluations (two) per iteration and thus can adjust the search direction more efficiently compared to DIRECT (either two or four). For a given limit on evaluation count, DITRI can more efficiently locate the optimal regions $\mathcal{P}_b$. This complies with Condition 3 for the proposed algorithm.

A.4. Evaluation point location

Point evaluation within a simplex is determined stochastically instead of (deterministically) at the center. The uniform sampling is such that the expected value for the evaluation point chosen is at the center of the simplex, $c = s_{n+1}$  

(A.5)

where $s_i = b_i$, $s_i = (0, 1, 2, ..., n_p + 1)$, and $b_i$ defines the coordinates of the $i$th simplex vertex.

The approach of a random evaluation point allows a faster convergence rate on average than that achieved by a deterministic evaluation. This is possible because the number of instances when a random sample is better placed than the midpoint is on average equal to the instant when it is worse placed. At the case when a random sample is worse placed, its effect would be dominated by the better placed sample and diminished at subsequent iterations. Finally, integer-valued design parameters are handled simply by shifting the design parameter $p_{i}$ to the nearest integer, or to the immediate larger integer if it is a half-integer, for any integer-valued design parameter $i$. This fulfills Condition 1b. The number of prediction steps $N$ is the relevant integer-valued design parameter.

The detailed outline of DITRI is given in Algorithm 2.

Algorithm 2. DITRI.

Require Design bounds $(p, \bar{p}),$

maximum evaluation $\bar{N}$ and iteration counts $\bar{N}$

1: $i = 0, i = 0$
2: Let $\Phi$ contain all evaluated points
3: for the two initial points $i = 1$ and $2$ do
4: Set normalized vertices of initial point $b_i$
5: Evaluate point $\mathcal{E}(p_i)$ in $b_i$ using (A.4), (A.5)
6: Record initial point in $\Phi$
7: $\overline{N} = \overline{N} + 1$
8: end for
9: $s_0 \leftarrow \{p_1, p_2\}$
10: loop
11: Find set $O$ of potentially optimal points in $\Phi$

Fig. A2. Simplex division over two iterations.
for all $\alpha \in \mathbb{O} \ do$

13. Divide simplex $b_1$ to obtain simplices $b_1^+; b_1^-$

14. For both $j = 1$ and $2$ do

15. Evaluate point $r(j)$ in $b_j$ using (A.4), (A.5)

16. Record new point in $\phi$

17. $i^{ev} \leftarrow i^{ev} + 1$

18. Update $S^{ev}$ to be all Pareto points in $\phi$

19. if $i^{ev} \geq n^{ev}$ then terminate algorithm

20. end for

21. end for

22. $i^* \leftarrow i^* + 1$

23. if $i^* \geq n^*$ then terminate algorithm

24. end loop

Appendix B. Proof of Theorem 26

Proof. Let $P(i^*)$ be the set containing all the evaluated points at iteration number $i^*$ in Algorithm 2. After some $i_0$ steps ahead, every point in $P(i^*)$ would eventually be classified as potentially optimal either via criterion (A.2) given its rank, or criterion (A.3) given that each search space size $d_j$ for all $j \in \{1, \ldots, iP(i^*)\}$ would become smaller than $d_j(i_0 + i_1)$. Each potential optimal search space will be divided and at least 1 new point will be evaluated after each division so that $P(i^* + i_1) \supset P(i^*)$. Consequently, as $i_0 \to \infty$, the algorithm would search the space $P_i$ entirely i.e. $\lim_{i_0 \to \infty} P(i^*) = P_i$. Any point $p \in P_i \subseteq P$ in the Pareto design set will be evaluated so that $\lim_{i_0 \to \infty} S_i = P_i$.

References


